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FORTRAN IV Compiler-Loader for the Wang 520/600 Calculator

March 1976



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Adelphi, Maryland 20783

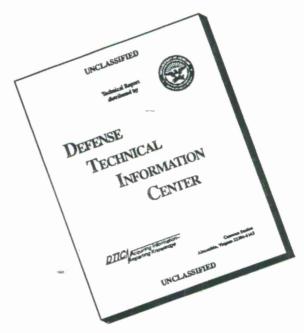
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# 20. ABSTRACT (Continue on reverse side if necessary and identify by block number)

This paper describes a compiler system that will translate FORTRAN IV programs into programs that will run on the Wang 520/600 programmable calculator. The system includes an option to generate the Wang programs on punched cards that can be input into the 520 or 600 via a mark-sense card reader.

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The compiler was written in FORTRAN IV for the IBM 370/195computer. Only a few minor changes are necessary to run the system on a different computer. The paper includes a detailed description on the types of optimization used to yield an effective FORTRAN translation onto a relatively small computer (i.e., desk calculator).

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#### 1. INTRODUCTION

In July 1973, a FORTRAN compiler system was announced that translated FORTRAN IV programs into Wang 520 calculator code. This system was developed on the Tym-Share Incorporated system. The output consisted of a listing of the translated Wang programs along with special tables needed by the user later for loading the program into the calculator.

A new compiler system is described here that contains the original system as a subset, but is written in FORTRAN IV instead of SUPER FORTRAN. The following characteristics describe the changes in the old system:

- (a) Loader--A loading phase has been added that will produce a binary card deck as an output of the system. This deck can then be read, via a mark-sense card reader, into the Wang 520 calculator.
- (b) Algorithm Changes--Several changes that have been made in the algorithms cited in the original paper either remove ambiguities or improve code translation.
- (c) New Optimization--Special attention has been placed upon the availability of the 16 basic registers to further optimize the code space. Also, many additions have been made to the second-pass optimizer, with special consideration to the optimization of array code generation.
- (d) Adaptability--Because the system is now written in FORTRAN IV, it should be easily adaptable to computers other than the IBM 360/370 series on which it is presently stored.

The following sections describe how the new improved system is used and give a more detailed account of the added characteristics.

## 2. USING THE SYSTEM

The compiler-loader system has been implemented on an IBM 370/195 computer, accessible at Harry Diamond Laboratories (HDL) through the IBM 1130 remote batch terminal. The following JCL statements are necessary to execute the Wang 520 compiler-loader system:

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973)

```
JOB CARD
ACCOUNT CARD

//JOBLIB DD DSN=C753.WANG,DISP=(OLD,KEEP)

// EXEC PGM=COMPILER,REGION=280K

//FT01F001 DD DSN=SCRATCH,DISP=(NEW,DELETE),UNIT=SASCR,

// SPACE=(1872,(133,5)),DCB=(RECFM=VS,LRECL=1868,BLKSIZE=1872,DSORG=DA)

//FT06F001 DD SYSOUT=A,DCB=(RECFM=FBA,LRECL=133,BLKSIZE=1596)

//FT07F001 DD SYSOUT=B

//FT05F001 DD *

Compiler-loader input

/*
```

The system output consists of the program source and object listings and special system tables, some of which have already been described. If the loader has been used, object programs are punched. However, since the terminal currently in use receives card images that contain only the 256 EBCDIC characters, the punched output is not in a form directly readable into the Wang 520. A program has been implemented on the IBM 1130 to convert the object decks to the Wang 520 card format that is acceptable to the mark-sense reader. The following deck should be submitted for local (IBM 1130) execution for obtaining the card conversion:

```
Col 1 4 8 11 51

// JOB 1210 your name

// BPNCH
```

cards from loader

blank Wang cards (same amount)

Each set of FORTRAN and Wang routines that should be loaded together is called a "load group." The punched output of the compiler-loader consists of a leader card with the number of cards (in hexadecimal format) in the object deck in columns 1 and 2. The object deck follows, with 40 steps punched on each card, also in hexadecimal format. The above sequence is repeated for each load group. Hence, in loading the compiler-loader system, sufficient blank cards should be added behind the system deck to take care of all the cards to be punched.

H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973).

The Wang cards punched correspond one to one to the cards punched by the loader, so the number of Wang cards required is the same as the number of cards put into the program. A maximum of 500 cards may be converted at one time.

# 2.1 Compiler-Loader Operating System

The compiler-loader represents a very simple operating system for which the user can run the program by using the set of five system operations (table I).

The general deck setup is

- (a) \*WANG card(s) for each external Wang program included in a load group,
- (b) \*FTC card(s) preceding each FORTRAN program deck included in a load group,
  - (c) \*LDR card(s) for load group,
  - (d) \*END.
  - (e) Repeat steps (a) to (d) for any additional load groups.
  - (f) \*STOP (always the last card in the deck).

FORTRAN Program Compilation.—Every FORTRAN main program or subroutine must be preceded by the \*FTC control card (see table I). The options NOLIST and NOMAP are used to limit the amount of printed output. The option AUTO is used to take advantage of the availability of free registers in a given program (see sect. 5.1.2(d)). The previous options need not be specified.

External Program Definition. -- The \*WANG card allows already existing Wang 520 programs to be combined with FORTRAN programs into one final loaded program. It is assumed that the name specified in the NAME field has been referenced by a FORTRAN program. All the options are used to avoid conflicts with marks and register numbers used by other programs included in the same load group. Hence, all marks and nonbasic registers used by the external program must be included. If the mark or nonbasic register range used by the program is not one definable sequence, the particular option can be repeated as often as necessary to describe all the sequences actually used. If more than one card is needed for a given external program, additional \*WANG cards can be added with the NAME field blank.

TABLE I. OPERATING SYSTEM COMMANDS

Oper	ation	Use	Program name required
(1)	FTC	FORTRAN program compilation	Yes
		Options:	
		(a) NOLISTsource program listing not printed	
		(b) NOMAPName table, translation, and reference table not printed	
		(c) AUTO (nn-mm)free registers in range nn-mm are assigned to most frequently referenced variables (if nn-mm is not specified (i.e., AUTO), range 01-15 is used)	
(2)	WANG	External program definition	Yes
		Options (all must be specified to prevent conflicts):	
		(a) Ennnnmark number of entry point to pro- gram (must be specified)	
		(b) Snnnnnumber of program steps, not including END PROG (if specified, storage is allocated for program before compiled programs)	
		(c) Mnnnn-mmmmMnnnn is mark number used by pro- gram (if -mmmm is specified, entire range is excluded by loader)	
		(d) Rnnn-mmmRnnn is nonbasic register number used by program (if-mmm is specified, entire range is excluded by loader)	
(3)	LDR	Assigns storage and marks to compiled programs and satisfies external references between programs	Optional
		Options:	
		<ul><li>(a) Mnnnnstarting mark number (if not specified, first available mark is used)</li></ul>	
		(b) Rnnnstarting nonbasic register number (if not specified, first available number is used)	
		(c) Ennnentry point mark number (if not specified, first available entry point number is used)	
(4)	END	End of load group	No
(5)	STOP	End of compiler input	No
		General Form	
Col	1	Col 2-5 Col 8-13 Col 16-71	
*		OP NAME OPT, OPT, , OP	Т

Program Loading.—The \*LDR card is used to assign storage and marks to compiler programs and satisfy external references between programs. If NAME is specified, it must have appeared previously on an \*FTC card. If NAME is blank, all programs compiled or declared external since the last \*END card are loaded. If one of a load group is named on an \*LDR card, each program must have a separate \*LDR card with the name used on the corresponding \*FTC card. None of the options need be specified.

End of a Load Group. -- The \*END card marks the end of a group of programs to be loaded together. Whenever the loader is used, an \*END card must follow the \*LDR card(s), to complete the loading process. There may be several load groups in one job, each terminated by an \*END card.

End of Compiler Input. -- The \*STOP card is used at the end of all compiler-loader input to halt execution of the system.

# 2.2 Control Card Examples

(a) Compile a single program and do not load. All listings are desired.

\*FTC MAIN
(FORTRAN deck)
\*END
\*STOP

(b) Compile and load subroutines A and B. Do not get listings of the translation; specify automatic register allocation for A. In the same job, compile and load subroutine C with external program D. Program D uses marks 0003 and 0010-0103 and registers 016-047; its entry point is 1010, and it uses 156 steps.

\*FTC A NOMAP, AUTO
(subroutine A deck)

\*FTC B NOMAP
(subroutine B deck)

\*LDR

\*END

\*WANG D E1010, S0156, M0003, M0010-0103, R016-047

\*FTC C

\*LDR

(subroutine C deck)

\*END

\*STOP

(c) Compile and load program INTG. Registers 03-07 are to be used for automatic register allocation. No source listing is desired. The loaded program uses marks beginning with 1200 and registers beginning with 032. The entry point for INTG is 1104.

\*FTC INTG AUTO(03-07), NOLIST (FORTRAN deck)
\*LDR M1200, R032, E1104

\*END

\*STOP

(d) Compile and load programs FTC and DERIV together. Program FTC should use marks beginning at 0800 and registers beginning at 025. Program DERIV uses marks starting at 0900 and registers beginning at 025. Even though the starting register number is the same, the loader does not assign the same registers to both programs. The first register used by DERIV is the first register after 25 not used by FCT.

\*FTC FCT

(FORTRAN deck)

\*FTC DERIV

(FORTRAN deck)

\*LDR FCT M0800, R025

\*LDR DERIV M0900, R025

\*END

\*STOP

# 2.3 FORTRAN Implementation and Limitation

The FORTRAN IV subset implemented by the Wang compiler-loader system has been completely described. To make this report as self-sufficient as possible, a general list of unimplemented FORTRAN IV capabilities is reported here (table II).

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973), section 2. Changes in the restrictions are noted in section 4 of TM-73-15S.

TABLE II. FORTRAN PROGRAM STATEMENTS REQUIRINT EDITING1

Subroutines not provided	Not implemented and ignored	Not implemented with error result	Must be followed by at least one blank
OVERFL(J) DVCHK(J) SSWTCH(I,J) SLITET(I) ERF(X) GAMMA(X) ALGAMA(X) CEXP(X) CLOG(X) CSIN(X) CCOS(X) CABS(X) CSORT(X) CMPLX(X1,X2) CONJG(X) AIMAG(X) REAL(X) DUMP(A1,B1,) PDUMP (A1,B1,)	ENDFILE PUNCH PRINT REWIND NAMELIST BACKSPACE FORMAT  Machine-dependent functions ignored AND(X,Y) OR(X,Y) COMPL(X) BOOL(X)	ASSIGN N TO I BLOCK DATA COMPLEX DATA ENTRY CALL EXIT EXTERNAL GOTO I, (N1,, NM) RETURN I(I not blank) Arithmetic State- ment Function	DIMENSION COMMON INTEGER REAL DOUBLE PRECISION LOGICAL SUBROUTINE CALL FUNCTION

#### Special cases:

- (a) Function names not allowed as arguments of a subroutine or function subprogram. Subscripted variables allowed when not outputs. Minimum of 1 and maximum of 15 arguments. Function nesting limited to 5.
- (b) Values of DO indices are not available outside of the loop.
- (c) Variables, constants, reserve words, or special operators cannot be continued on the next line.

In general, all the important facilities in FORTRAN IV have been implemented; perhaps the only two deficiencies are no complex arithmetic and only single-dimension arrays. However, these restrictions are sensible when the computer being used has the limited size of a Wang 520. Table III summarizes the specific restrictions.

<sup>1</sup> From H. Bloom and A. Hausner, Harry Diamond Laboratories TM-73-15.

# TABLE III. SET OF FORTRAN IV CAPABILITY RESTRICTIONS

- (1) Complex constants are not allowed.
- (2) Arrays can be only one dimensional.
- (3) Trigonometric functions arguments such as SIN have a magnitude restriction of 10 radians.
- (4) Computed GOTO has a limit of 20 statement numbers. If the index is out of range, the default is to the last number.
- (5) The DO index value is not available outside the DO loop. The maximum DO nest level is 5.
- (6) The "n" in PAUSE or STOP may be only a single decimal digit.
- (7) Maximum and minimum built-in functions can have two arguments only.
- (8) User-defined subprograms do not supersede system routine names.
- (9) Variable dimensions are not allowed.
- (10) Only one labeled COMMON can appear on a given COMMON card.

The compiler-loader system assumes that the FORTRAN decks submitted are free of syntax errors and thus does not perform any degree of syntax checking itself. Therefore, it is possible for a run to fail, and the computer will list an error that is completely meaningless to the user. In this case, the user should satisfy himself that his decks are running programs.

### 2.4 User-Supplied Compiler Commands

As described in detail previously, the system allows the user to include special compiler commands (table IV) to improve his code, if he desires. These commands appear directly in the FORTRAN deck and are entered as special comments. The commands T, R, and S are used for register optimization. Since the automatic register assigner subsystem (AUTO, sect. 3.3) was added, there is little need for the user to assign variables to basic registers. However, the option is available and does

<sup>&</sup>lt;sup>1</sup> Summarized from H. Bloom and A. Hausner, Harry Diamond Laboratories TM-73-15, section 2.

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973).

TABLE IV. USER-SUPPLIED COMPILER COMMANDS

Command	Effect
1. T Z (Top assigning)	The top register in the program is changed from 15 to Z ( $<$ 15).
2. R N(V <sub>1</sub> ,,V <sub>m</sub> ) (Equivalencing)	The FORTRAN variables $, V_m$ are all assigned the same basic register $N(00 \le N \le 15)$ .
<ol> <li>S name(C<sub>0</sub>,,C<sub>m</sub>),</li> <li>d, R1-R2, \$Z</li> <li>(Subprogram</li> </ol>	(a) "name" is the name of a subprogram to be called as name(a <sub>0</sub> ,,a <sub>m</sub> ).
specification)	(b) C <sub>i</sub> is a special symbol that determines that property of the corresponding actual subprogram call argument as follows:
	C =B; argument a is both an input and output variable (a must be nonsubscripted).
	I; argument a is input expression only.
	O; (letter O) argument a is output variable only. (ai must be nonsubscripted).
	E; argument a is empty variable, i.e., neither input nor output.
	(c) d is the maximum number of nested DO indices in "name." If d is omitted, all necessary DO indices of the calling program will be saved and restored.
	(d) R1-R2 is the range of basic registers to be saved in the calling program. If R1=R2=0, no registers are saved. If the range is omitted, all variables assigned to basic registers in the calling program will be saved and restored.
	(e) Z is the top register specified in a T statement in "name." If \$Z is omitted, Z=15.
	(f) d, R1-R2, and \$Z fields may be listed in any order.

<sup>&</sup>lt;sup>1</sup>From H. Bloom and A. Hausner, Harry Diamond Laboratories TM-73-15 (July 1973)

TABLE IV. USER-SUPPLIED COMPILER COMMANDS (CONT'D)

Command	Effect
4. W V(format, N, M) (I/O specification)	(a) V is the name of a FORTRAN variable.
	(b) Format is any four digit allowable Wang format that can follow the PRINT step (table 33 of HDL-TM-73-15).
	(c) N is the print-on-read indicator.
	N = 0; do not print input on READ. $N = 1$ ; print input on READ.
	(d) M is the spacing indicator.
	<pre>M = 0; no spaces. M = 1; space before printout. M = 2; space after printout. M = 3; space before and after printout.</pre>
	(e) If N = 0, M is ignored on a READ.
	(f) If format is 0015, N and M are used with the previously defined default format (table 19 of HDL-TM-73-15).

- (a) All commands must start with COMPILE in columns 1-7.
- (b) Card commands cannot be continued on next line.
- (c) Only one S command can appear for each "name."
- (d) Any FORTRAN variable can appear in at most one R command.
- (e) Subprogram arguments cannot appear in an R command.
- (f) Only one register can be specified in one R command.

<sup>&</sup>lt;sup>1</sup>From H. Bloom and A. Hausner, Harry Diamond Laboratories TM-73-15 (July 1973)

allow for the most efficiently generated code. The command T is usually used with respect to the command S, in order to line up arguments in the called routine with the argument register storage in the calling routine. The S command is mainly designed to minimize the amount of code that must be generated each time a call is translated. Examples of ways that the R, S, and T commands can be used have been given earlier.

Since format statements are not implemented in this system, the W command allows the user to make full use of the I/O capabilities of the Wang 520 printer. Hence, whenever a variable is referenced in an I/O statement, the indicated W format is used in the translation, if it exists.

#### 3. DESCRIPTION OF SYSTEM STRUCTURE

The compiler is divided into six subsystems: (1) encoder, (2) parser, (3) automatic register assigner, (4) translator, (5) optimizer, and (6) loader. A brief description of each system is given below.

## 3.1 Encoder

The encoder takes the FORTRAN program string code and generates an encoded program string. It performs the following special features:

- (a) Generates special operator code (for example,  $X^2$ , 1/X, II, power of 10 shifts) to take advantage of the Wang operator set.
  - (b) Sets up COMPILE option tables.
- (c) Eliminates FORTRAN formats from code and removes format number and device index from I/O statement.
  - (d) Removes commands not allowed in version.
- (e) Generates special end-of-DO (CONTINUE) statement with new label that also replaces the DO statement label. This is a new feature that eliminates ambiguities in the pretranslation of a FORTRAN statement into a multiline code (i.e., for I/O, certain LOGICAL IF and function calls).

## 3.2 Parser

The parser takes the encoded string and generates a reverse Polish string. It performs the following special features:

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973), table 12, section 3.2.1.

- (a) Reorders expressions having a commutative binary operator (+, \*, .AND., .EQ., .OR., .NE.), if there is a reduction in the number of registers needed to compute the expression. The new algorithm replaces the one mentioned previously  $^{\rm l}$  in that it simply checks the expression with the arguments in either position to determine the best ordering.
- (b) Reorders commutative expressions, provided there is no change in the number of registers under the following conditions (in order of priority):
- (1) Reorders the assignment expression so that a variable assigned in the previous statement is closest to the left of the present expression. Doing so allows the recall suppression feature to be more effective; for example, statements k = 1, I = J + K will be pretranslated as K = 1, I = K + J.
- (2) Reorders the expression to replace the two consecutive operators "- $_{unary}$ " + with "- $_{binary}$ "; for example, X = -A + B becomes X = B A.
- (c) Defines code for (F, (A, and (M to indicate the end of a function, array, and macro argument list, respectively.
- (d) Checks LOGICAL IF statement for function or I/O statement expansion candidates. If expansion is required, the statement is altered. For example,

IF(X. EQ.1) Y(I) = F(X)

becomes

IF (X. EQ.1) GOTO L1

GO TO L2

L1 Y(I) = F(X)

L2 next statement

where

L1 and L2 are system-generated labels.

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973), table 12, section 3.2.1.

- (e) Stores equivalence and common statements into system tables.
- (f) Checks for function or I/O expansion. If the statement has a label, the label goes with the first statement in the expansion. For example,

10 
$$Y(I) = F(X)$$
 becomes 10  $T.00 = F(X)$ 

$$Y(I) = T.00$$

(g) Defines the array repeat operator  $A_{\rm R}$  if assignment array appears on the right side of the assignment. For example,

$$X(I) = X(I) + B/X(I)$$

becomes

$$X(I) = A_R + B/A_R$$

- (h) Looks for the pattern "expression \* expression" (appearing in FORTRAN format) and replaces it with (expression)\*\*2.
- (i) Looks for the pattern ++ or \*\* (appearing in reverse Polish). If it is found, the second operand of the first operator is switched with the operator. For example,

the FORTRAN (A + B) + (C\* D), first in Polish (AB+CD\*+), and then (CD\*AB++) after reordering becomes

the FORTRAN ((C \* D) + A) + B, or in Polish (CD\*A+B+).

# 3.3 Automatic Register Assigner (AUTO)

The AUTO system assigns auxiliary registers (i.e., work registers, DO index registers, and register 00) to common subexpressions and temporary variables and free registers to other variables if so specified. The discussion of register optimization is very important and is handled specially in section 5.1. This system is new to the translator, but the system can greatly reduce the program size by trying to optimize the use of the basic registers.

## 3.4 Translator

The translator takes the parsed string and generates the relocatable Wang code. There are many local optimization features performed at this point.  $^{\rm l}$ 

The techniques used in translating each type of FORTRAN statement into corresponding Wang Code have been described in detail. Even if the compiler-loader system is not used, it is worthwhile to learn the translation algorithms and apply them when writing the Wang code, especially in the case of conditional transfers and DO loops. The original report also describes the basic translator scheme for setting up the calculator as a computer structure and use of the registers in implementing the translation algorithms.

# 3.5 Optimizer

The optimizer (or second-pass optimizer) takes the relocatable Wang code generated by the translator and searches for certain patterns, in order to reduce the number of steps. The seven passes made through the code have also been discussed. The current version of the optimizer makes an additional eight passes to try to further reduce the number of steps. These are fully discussed in section 5.2.

# 3.6 Loader

The loader combines the whole relocatable Wang code generated from the FORTRAN programs with any externally defined Wang programs and stores the whole code together in absolute form by use of the following order:

- (a) External Wang programs
- (b) FORTRAN programs
- (c) Large number storage in nonbasic registers
- (d) Nonbasic variable register storage.

The loader outputs a set of punched cards representing the object deck of a loaded program module (sect. 2).

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973), tables 15, 17, 18.

An optimization algorithm within the loader attempts to use the same nonbasic registers for variables in different programs, if doing so is possible. Variables (including arrays) that are candidates for sharing a nonbasic register with a variable (or array) in another program must obey the following property. Between the first and last reference to the variable, there can be no call to another subprogram. This restriction is necessary to keep the value of the assigned nonbasic register from changing in the called subprogram and, hence, storing an incorrect value for the variable used in the calling program.

#### 4. CHANGES IN ALGORITHMS

Aside from many new features in the compiler system, there have been many changes also in the algorithms already developed and described.  $^{\rm l}$ 

(a) The code generated from the nonbasic operators .LE., .LT., .GE., and .GT. has been shortened, and the code for the .EQ. and .NE. simple LOGICAL IF statement has been changed, to avoid the use of register 00. Table V shows the changes.

TABLE V. CHANGES IN LOGICAL CODE

.LE.	.LT.	.GE.	.GT.	IF (J.EQ.K) S	IF (J.NE.K) S
- L	- L	- L	- L	recall J	recall J
T L	T L	E 1	E 1 .	ST L	ST L
J IF O	J IF O	SP-ST L	SP-ST L	recall K	recall K
E 0	J IF +	J IF O	J IF O	- L	- L
J IF +	E 1	J IF +	J IF +	J NE O	J IF O
E 1	ST L	T L	GO	translate S	translate S
ST L	RE L	GO	T L		
RE L		RE L	RE L		

(b) The code used for generating function expansion has been altered so that if a label appears in the original statement, this label eventually appears in the first statement of the expansion.

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973).

- (c) The code for generating DO loops has been altered. The special DO label Lnn is now used as the label on the last statement of the DO loop. (A special CONTINUE statement is added, see sect. 3.1(e).) The combination of changes (b) and (c) removed the problem of not allowing function calls to be in the last statement of a DO loop.
  - (d) The special DO label name generated is now D.nm.
  - (e) The system label generated is now LABnm.
  - (f) The system-created variable (ZZZZnm) is now T.nm.
  - (g) The subprogram save variable (AAAAnm) is now AAA.nm.

#### 5. NEW OPTIMIZATION FEATURES

Since the development of the original compiler, the major change in the new system (besides the loader phase) has been the concern for more step optimization. The work lies in two major areas:

- (a) maximum use of the basic registers
- (b) recognition of repeatable patterns in the code so that copies of the patterns can be eliminated (e.g., generation of the same array element in two statements). The two types of optimization take place in the AUTO and optimizer subsystems, respectively.

The user previously had to decide whether or not he needed optimization tricks. They are now automatically performed by the compiler (table  ${\tt VI}$ ).

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973).

TABLE VI. ADOPTION OF OPTIMIZATION TRICKS

Туре	Section with reference 1
Using the basic registers effectively	5.1 (6.2.1)
Arithemetic IF statements	5.2 (6.2.2)
DO index suppression	3.2-b-1 (6.2.4-2)
Low index (of DO) is an expression	5.1.2-b Case 2 (6.2.4-3)
Register 00 is not used (upper index is expression)	5.1.2-b Case 3 (6.2.4-4)
Recall array suppression	5.1.2-a (6.2.5-1)
Transposing terms for recall suppression	3.2-b-1 (6.2.6-1)
Transposing terms to lower work register requirements	3.2-a, 3.2-i (6.2.6-3)
Trailing zeros	5.2 (6.2.7-1)
Constant set	5.2 (6.2.7-2)
Decimal suppression	5.2 (6.2.7-4)
Squaring	3.2-h (6.2.8-3)
ZZZZnn equivalencing	5.1.2-c (6.2.8-4)

<sup>1</sup> Section number in parentheses is from HDL-TM-73-15.

# 5.1 Register Optimization

As mentioned in section 3.3, the most important resources in the calculator are the 16 basic registers that play an important role in the statement translation. Some of the optimization concepts were obtained from Gries, <sup>2</sup> Hopgood, <sup>3</sup> and Rustin. <sup>4</sup> However, it is important also that some of the registers be used for storing the values of variables. There are three very critical savings:

<sup>&</sup>lt;sup>2</sup>D. Gries, Compiler Construction for Digital Computers, New York: John Wiley and Sons (1971).

<sup>&</sup>lt;sup>3</sup>F. R. A. Hopgood, Compiling Techniques, New York: American Elsevier Publishing Co. (1969).

<sup>&</sup>lt;sup>4</sup>R. Rustin, ed., Design and Optimization of Compilers, Englewood Cliffs, NJ: Prentice Hall, Inc. (1972).

- (a) Since "nonbasic" registers are constructed by the concatenation of eight program steps, these steps can be saved if the variables can be stored in basic registers, instead.
- (b) Many of the step optimization features (such as simple updating) can be applied only if the varible is in a basic register.
- (c) Execution time is saved, not only from the above types of step savings, but also by the use of basic register store-recall rather than "nonbasic" store-recall.

The goal of register optimization is to make use of the basic registers when they are not being used in the ordinary statement translation. Let

 $\mathbf{M} = \mathbf{the} \ \mathbf{maximum} \ \mathbf{number} \ \mathbf{of} \ \mathbf{work} \ \mathbf{registers} \ \mathbf{used} \ \mathbf{in} \ \mathbf{any} \ \mathbf{one} \ \mathbf{statement}, \ \mathbf{W}$ 

 ${\rm M}_{\rm d}$  = the maximum number of DO index registers needed at any time,

LT = the first available work register,

W, = the number of work registers needed for statement i,

D; = the DO index level at statement i.

For statement i, the registers LT-W through LT-M +1 and Md through D +1 are free for some other use. In addition, register 00 might also be free. Hence, these registers can be used for storing the values of variables at selected points in the program. In programs where there is a large value for W (such as 4) or D (such as 3), register optimization can result in tremendous savings of space, since these peak values occur infrequently in the program (for W , usually once or twice). However, care must be taken in the choice of which variables to use with which registers. The next section lists definitions of terms needed in the description of the register optimization that is performed by the compiler.

## 5.1.1 Definitions

(a) Auxilliary register—any register that at any time was used as either a work register (i.e., for computing expressions) or a DO index register. Register 00 is also considered auxiliary. These registers are used at least once in the ordinary statement translation described in section 3.4. An auxiliary register is considered available in a given statement span if it has not been previously used in any statement within the span.

- (b) Free register--free registers F are those defined by LT-M  $\geq$  F > M. These registers are never used in the ordinary statement translation; hence, they are entirely free to be used for variable register assignment.
- (c) Variable domain--the span of statements from the first to the last referencing the variable in the program. If the variable appears outside a DO loop and its last reference is within the DO loop, its domain is extended to include the entire DO loop.
- (d) Common subexpression domain—the span of statements from the first to the last that includes the subexpression. Within the span, the following criteria must be satisfied:
- (1) No statement labels except for the following special case:

IF (arithmetic expression) Ll, L2, L3

 $L_{i}$  statement where i = 1, 2, or 3.

(This criterion is somewhat restrictive, but greatly simplifies the pattern searching. Also, it eliminates the need to move the expression backwards in the program to make certain that no jump can occur over the first appearance of the expression.)

- (2) No new assignment involving a variable used in the expression
  - (3) No subprogram references.
- (e) Closed block--a span of statements obeying the following properties:
  - (1) No backward transfers
- (2) No transfers from outside the block to within the block
  - (3) DO loops may appear within a closed block
- (4) No function or subroutine call can appear except as the first statement of the closed block.

(Without properties (1) and (2), the span of the closed block would be inefficiently expanded to include the whole range from the transfer back to the label reference. Doing so would eliminate most of the register optimization. Property (3) is included because of the popular usage of the DO loop, which usually dominates every segment of any program. Property (4) is necessary because the compiler does not know what registers may be used in the subprogram; hence, it would have to save all registers used up to that point and result in extra steps generated by the save code.)

- (f) Local domain of a variable--any span of statements that references the variable and obeys the following properties:
- (1) The first statement is an assignment to the variable (not as a result of a LOGICAL IF). This cannot be an update (Y = f(Y)).
  - (2) All properties of a closed block must be obeyed.
- (3) The span includes all references after the first assignment that obey properties (1) and (2).

(The properties define a region in which a variable is totally defined, since the first statement must be an absolute assignment.)

(g) Temporary variable—a variable whose domain lies entirely within a single closed block. In addition, the variable may not have appeared in a COMMON, FUNCTION, SUBROUTINE, or EQUIVALENCE statement, nor be an array name. A good example of a temporary variable is the switch variable used for arrays:

$$TEMP = A(I)$$
,  $A(I) = A(I+1)$ ,  $A(I+1) = TEMP$ .

The variable TEMP has no real importance, except to be used as temporary storage. Certainly, TEMP should be stored in a basic register, if possible.

- (h) Global variable—a variable whose domain includes one or more closed blocks. These variables are usually heavily referenced throughout the program; hence, they cannot fit within any defined domain. However, it would be especially beneficial if basic registers could be assigned them because of the step savings in the referencing.
- (i) Local variable——A variable over whose domain exists one or more local domains. The local variable obeys the properties of a temporary variable. Local variables are, in most cases, special examples of temporary variables. In general, the programmer has used

the name of a variable that is constantly used for temporary storage in many places in the program. The local domain definition implies that the variable is entirely defined within a small local region, just as in the example given for the temporary variable——see (g) above. This local definition may occur several times for the same variable in one program.

# 5.1.2 Register Assignment

As stated in section 3.3, auxiliary registers are assigned by the system in its process of translating the FORTRAN statements. This section is concerned with the assignment of registers for variables. The system does assign registers to all arguments of a subprogram (fig. 1). It can allow the programmer to specifically assign variables to registers under program control (sect. 2.4). The system register assignments are made in the following order:

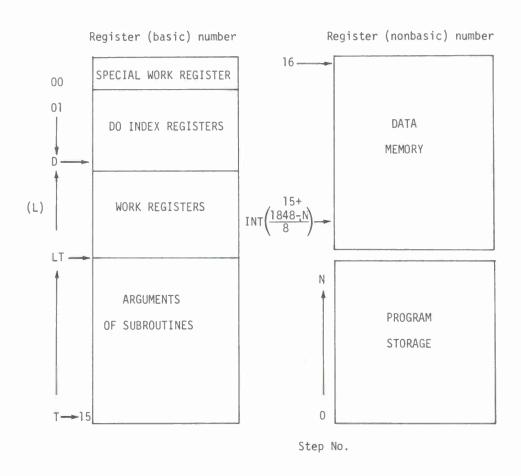


Figure 1. Configuration of translator scheme.

(a) Assignment of common subexpressions—If auxiliary registers are available, common subexpressions are assigned to them in the following order. Let p = the number of occurrences times the number of tokens (e.g., operators and operands) in an expression. Hence, p is a rough estimate of the number of steps that can be saved by replacing all occurrences of a common subexpression with a variable reference and initially assigning the variable the value of the expression. The expressions are considered in decreasing values of p. The register used is no longer available in the entire expression domain. If a common subexpression cannot be assigned a register, it is ignored.

Common subexpression register assignment was chosen first because the domain is so restrictive and huge savings of steps are possible. Experience has shown that most subexpressions are array elements (such as A(I+1)). In this compiler, it takes at least six steps to generate an array element. Hence, if the element appears three times, at least 9 steps can be saved (12 steps, less two extra recalls and one store). Expressions are ignored if there are no free auxiliary registers, since use of a temporary variable for storage requires the generation of a nonbasic register that requires an additional eight steps. Experience has shown, however, that a free auxiliary register is almost always available.

(b) Assignment of a temporary variable—temporary variables are assigned to auxiliary registers, if available, in the order of decreasing frequency of references. The register used is no longer available in the entire variable domain.

Approximately 20 percent of the variables in an average program obey the definition of a temporary variable. However, only a few of these variables have an auxiliary register available over their domain, unless the domain is relatively small in its span.

In three special cases, the auxiliary register availability definition can be overridden and provide even more optimization. The cases follow:

Properties of Variables Considered Special Cases

- (1) Exactly two references are in consecutive statements (in the entire program).
  - (2) First reference is an assignment.
  - (3) Second reference is not in a labelled statement.

### Cases

(1) Strictly temporary variable:

X=. . . ) Register LT is used for X if A is nonarray and LT-1 is used if A is an array.

or

For this case, X (or K) would have been recalled and then stored in LT anyway; instead, the values can be stored in LT directly, even though LT is not available according to the definition. Not only is a register saved, but also steps are saved.

(2) Low limit of DO loop:

Case (2) occurs when the lower limit is an expression. Since FORTRAN IV does not allow expressions as DO limits, the programmer must resort to the use of a temporary variable. The system essentially undoes this wasteful operation.

(3) Upper limit of DO loop:

The system uses register 00 to store the contents of the upper limit at the end of the DO loop in preparation for a transfer check back to the beginning of the loop. As in case (2), the system essentially allows the upper limit to be an expression.

(c) Assignment of local variables—local variables are assigned to auxiliary registers, if available, in the following order (each is based on decreasing frequency of references): (1) variables whose every reference is included within a local domain and (2) variables for which at least one reference does not appear in a local domain.

A distinction must be made between cases (C-1) and (C-2) because the variable in case (C-1) is always redefined at the beginning of a local domain. Hence, it makes no difference if the auxiliary register used to store the variable value is reused for another purpose between two local domains of the same variables. However, for case (C-2), it is necessary to add a special store instruction each time an assignment is made, so that a global value can be maintained for the variable.

It is possible that a variable can be defined as a local variable in many local domains. Hence, in one domain, the variable could be assigned register 10 and in another domain, register 2. The local variable register assignment is perhaps the most heavily used portion of the entire register assignment algorithm, since the local domains appear in such frequency and are, in general, extremely restrictive in domain size. Cases (b-1), (b-2), and (b-3) for temporary variables apply also for local variables. The register used is no longer available in the entire local domain.

The following example should help to illustrate the concept of local variable register assignment: Assume that X is assigned registers 1 and 2 in each domain, Y is assigned register 2 in its first domain, and Z is assigned register 8.

Code

X = 1E 1 Y = X+1ST 1 E2 Z = Y+2RE 1 E4 ST2 STORE ST 2 GO TO 20 El +2 MARK E2 20 STORE Y = 4Y ST LT Y = 2+YST 8 RECALL E 2 Y +8  $\times$ LT X = 2SEARCH STORE 20 Y

The variable X satisfies case (c-1) since it is totally defined within every domain where it appears. Hence, a register can be simply substituted for X for each reference. However, Y appears as case (c-2). Its first domain is local, but not the remaining references. The programmer can substitute a register for Y in the first domain, but this register may not be available at statement 20. Hence, one must also save the value for Y in its general storage location.

(d) Assignment of free register variable—the programmer can specify that the free registers be used for variable assignment (AUTO option, table I). They are allocated in order of decreasing frequency of references to all variables not already assigned registers by the system. Only one variable can be assigned to a given free register.

Since the compiler is designed to store arguments of subroutines in basic registers, care must be taken in subprogram calls to save the values of variables used for permanent storage in nonbasic registers if they would be destroyed by the called program. Hence, the best feature of the common-expression, temporary variable, and local variable register assignment is that the values in the registers can be destroyed when calls are made. However, if free registers are assigned, their contents may have to be saved. The programmer does have control in specifying which registers need be saved when the program calls a given subprogram (command S, table IV).

# 5.1.3 Example: Use of Auxiliary Registers

The sample program of figure 2 illustrates the use of auxiliary registers. A register map is given in figure 3. The ARG column gives the assigned register number (20 indicates that a nonbasic register must be assigned). A maximum of three work registers and one DO index register is used. (Actually, there are only two registers after B(I) is set up as a common subexpression. See A(I) statement.) this example, all the variables qualify as temporary variables, and the variables E, F, and X are assigned registers 00, 14, respectively. The variables C, M, and K are strictly temporary variables that can be assigned register 15. The temporary variable N satisfies the lower DO index condition and hence is assigned register 01. The temporary variable KK satisfies the upper DO index condition and is assigned register 00. The variables A and B are

assigned nonbasic registers. The variable J, although temporary, cannot be assigned an auxiliary register over its domain, since none is available. However, its domain can be subdivided into two local domains, each of which satisfies the strictly temporary variable case (J is assigned register 15). Since, in general, J can be assigned a different register in each domain, its ARG column has the value 20.

#### FORTRAN COMPILER TESTI

```
C
      PROGRAM ILLUSTRATES THE USE OF AUXILIARY REGISTERS
      DIMENSION A(1), B(1)
      F⇒1
      F=1
      C = 1
      X=F+F*C
      M=2
      J=1+M
      K=2*J
      N=K+1
      DO 10 I=N.5
      A(1)=B(1)+1,/(B(1)*B(1))
   10 CONTINUE
      KK=B(I)
      DO 20 I=1.KK
      X = X + 1
   20 CONTINUE
      J=3
      X=J*J
      STOP
      END
```

Figure 2. Sample program listing.

#### NAME TABLE

INDEX	NAME	DIM	TYPE	ARG
1	Α	1	0	20
2	В	1	0	20
3	Ε	0	0	0
4	F	0	0	14
5	С	0	0	15
6	Χ	0	0	13
7	М	0	1	15
8	J	0	1	20
9	K	0	1	15
10	N	0	1	1
11		0	1	20
12	KK	0	1	0

Figure 3. Register table.

It is assumed that the loader will store the pointer value n-1 in the register allocated to A if register n contains A(1).

# 5.1.4 Program Execution

The relative Wang code generated is shown in figure 4. The program illustrates the following points:

- (a) Remembering constants--Steps 0 to 3 store 1 into E, F, and C. However, since X has been assigned register 13, the system changes ST 15 to St 13 to save a step and computes X in register 13.
- (b) Expression optimization—Steps 4 to 6 translate X=E+F\*C. The parse string is XCF\*E+=. Since the system remembers that F is in the window, it does not recall it, but simply multiplies the value by the result of X obtained so far (step 4).
- (c) Strictly temporary variables--Steps 7 to 10 translate M=2; J=1+M. Steps 11 to 13 translate K=2\*J; N=K+1. Even though M, J, and K are using register 15, the register number is changed to 01 to correspond to the first assignment for N. Step 11 illustrates the use of the special "2\*" operator that does a repeated addition to save a step.

STEP	CODE	BUTTON	KEY	S	TEP	CODE	BUTTO	N KEY
0	0001		E 1 00 14 13 13 00 13 E 2 01 E 1 01 01 E 1 01 MARK	35	00	05		E 5
1	0600	ST	00	36	06	00	ST	00
2	0614	ST	14	37	00	01		E 1
3	0613	ST	13	38	02	01	+	01
4	0413	Χ	13	39	09	02 S	HIFT	ALPHA
5	0700	RE	00	40	08	06		SIN
6	0213	+	13	41	08	00		SEARCH
7	0002		E 2	42		D00		
8	0001	ST	01	43	00	01		E 1
9	0001		E 1	44	06	15	ST	15
10	0201	+	01	45	08	01		RECALL
11	0201	+	01	46		В		
12	0001		E 1	47	02	15	+	15
13	0201	+	01	48	15	11		INDIR
14	0900	SHIFT	MARK	49	07	15	RE	15
15		D00	15 RECALL	50	09	12 S	HIFT	INT
16	0615	ST	15	51	06	00	ST	00
17	0801		RECALL	52	00	01		E 1
18	0035	В	15 INDIR 15 00 01 15	53	06	01	ST	01
19	0215	+	15	54	09	00 S	HIFT	MARK
20	1511	DE	INDIR	55	0.0	100		
21	0/15	KE	15	56	00	01		E I
22	0600	21	00	5/	02	13	+	13
23	0/01	KE	01	58	00	01		E I
24	0001	21	DECALI	59	02	01	+	01
25	0801	Λ.	RECALL	60	09	02 5	HIFI	ALPHA
27	0215	A .	16	01	00	00		21N
20	0213	DE	00	62	UO	נחח		SEARCH
20	0614	CT	1.4	6.4	00	กว	4	ГЭ
50	0014	31	V**2	65	06	15	СТ	L 3
3U	0012		1/4	66	00	10	21	V**3
31	0013	+	1/ ^	67	06	12	СТ	12
33	1511	1	TNDTD	60	00	US CI	HIET	STUD
34	0615	ST	15 00 14 X**2 1/X 14 INDIR 15	60	Na	14 9	HIFT	FND
34	0013	21	10	09	U9	14 31	HILI	LIND

Figure 4. Wang code listing.

- (d) The DO loop label--Steps 14 to 15 define the DO loop DOO. The system knows that the index N is immediately available and hence can store it in the first available register used by computing  $B\left(I\right)$ .
- (e) Common subexpression--Steps 16 to 22 are used to compute the common subexpression B(I) and store it in 00.
  - (f) Squaring--Step 30 is used to find B(I)\*B(I).
- (g) End-of-DO-loop--Steps 35 to 42 perform the end-of-DO loop check ( $\alpha$ , SIN causes a two-step jump if the contents of register 00 are less than the contents of the display window).
- (h) Upper DO index--Steps 43 to 51 translate KK = B(I), storing the result in register 00. Doing so saves the updating of register 00 at the end of the DO loop (see steps 35 to 36 for the previous DO loop).

In many test cases, as much as 20 percent of the core was saved due to the register optimization.

# 5.2 Additions to the Optimizer Subsystem

The second-pass optimizer (sect. 3.5) contains eight additional passes through the translation code in an effort to reduce the number of steps:

(a) This pass effectively changes the FORTRAN statement,

IF(A)N1,N1,N2 to IF(A.LE.O.)GOTO N1

GOTO N2

unless statement N1 follows. In that case, IF(A.GT.O.) GOTO N2 is used. Also, the statement

IF (A) N1, N2, N2

is changed to

IF(-A.GT.O.)GOTO N1 if statement N2 follows

or

IF (-A.LE.O.) GOTO N2 if statement N1 follows.

If neither follows, no change is made.

(b) This pass changes the sequence

Jif+		ch sign
search		Jif+
Ll		search
search	to	L2
L2		Mark
Mark		Ll
Ll		

It does this before the pass that eliminates unreferenced marks, so that Mark, L1 is eliminated also if it is not referenced. The code arises from the FORTRAN statement IF(A), L1, L3, L2 with statement L1 following.

- (c) This pass checks to eliminate trailing zero digits in a number when advantageous.
- (1) If there is only one trailing zero of an integer, no change is made.
- (2) Two or more trailing zeros of an integer are changed to exponential notation; for example, 1000 is changed to 1.E3 with the translated code E1,  $\alpha$ , f3.
- (3) Trailing zeros of decimal numbers are eliminated; for example, 12.7300 is changed to 12.73.
- (4) If the first trailing zero of a decimal number is left of the decimal point, the decimal point is eliminated, and steps (1) and (2) apply; for example, 12500.00 is changed to 125E2, with the  $\alpha$  shift used to multiply by 100.
- (5) All the trailing zeros of a number with an exponent to the right of a decimal point are eliminated; for example, 12.50E23 is changed to 12.5E23. This change does not result in an optimum translation, since the decimal point can be eliminated by use of 125E22. It is eliminated in pass (d).
- (6) If the first trailing zero of a number with an exponent is to the left of a decimal point or just to the right of it, the decimal point is eliminated, and the exponent is adjusted; for example, -120.00E25 is changed to -12E26, and 32.0E-15 is changed to 32E-15.
  - (7) The code EO, ST R, is changed to T R.

The net effect of this pass is to make sure that no number is written with a trailing zero (in the mantissa, if in exponential notation) unless it is an integer multiple of 10, but not of  $10^n$  for  $n \ge 2$ .

- (d) This pass insures that identical numbers have identical representations.
  - (1) Integers are left untouched.
- (2) Decimal points are introduced, if an exponent can be eliminated; for example, 12E-1 is best written 1.2.
- (3) An exponent of 1 is eliminated for a nondecimal mantissa; for example, 12El is changed to 120.
- (4) A SETEXP notation is changed to an  $\alpha$  shift by introduction of a decimal point; for example, 12E-16 is best written 1.2E-15. The exponent field can now be translated in two steps  $(\alpha, F15)$  instead of four, with a net saving of one step.
- (5) All leading zeros of decimal numbers are eliminated; for example, 0.03 is represented as 3E-2 and 0.0001 as 1E-4.
- (6) For a decimal point with an exponent, the exponent is eliminated, if possible, by shifting the decimal point to left; for example, 12.3E-2 is changed to .123.
- (7) Otherwise, a check is made to see if a SETEXP notation can be changed to an  $\alpha$  shift; for example 12.3E-16 is best written 1.23E-15 as in item (4).
- (8) Otherwise, a check is made to see if the exponent can be eliminated by moving the decimal point to the right; for example, 12.36E2 is changed to 1236, and 0.013E2 is changed to 1.3.
- (9) Otherwise, a check is made to see if the decimal point can be eliminated by moving it to the right; for example, 12.3E5 is changed to 123E4 (but not 12.3E-15 to 123E-16). See item (4) about A SETEXP notation.

The net effect of passes (c) and (d) is to obtain a unique translated representation for all numbers.

(e) This pass eliminates the necessity of entering a number when it is determined that it is in the window. Thus, the FORTRAN code

A(1)=1.2

X=1.2

Y=1.2

is translated

generated [A(1)] in LT
E 1
E .
E 2
INDIR
ST LT
STORE X
STORE V

- (f) This pass further attempts to optimize steps involving arrays.
- (1) If the value of a needed array is in the window, it is not generated again. For example,

$$A(I) = \dots$$

$$T = A(I) + \dots$$

is obvious.

(2) If the index of an array is determined to be in a register, it is not generated again. For example,

$$A(I) = ...$$
  
 $X = Y$   
 $T = A(I) + ...$ 

The determination is made that  $\ell[A(I)]$  is still in LT, so the second A(I) is translated INDIR, RE LT. Another example is

where T is a basic register. The computation is performed in that register, leaving  $\ell[A(I)]$  and  $\ell[B(I)]$  undestroyed.

(3) Item (2) applies if the index differs only by a constant. A code to subtract or add the constant is supplied as needed. This pass was motivated by frequent use of an exchange code in problems involving arrays. Thus, the code

is now translated with the locations of A(I) and B(I) generated only once.

- (g) This pass tries to avoid an extraneous code involving the use of register 00.
  - (1) It eliminates T 00 when not needed. For example,

IF(A)5,5,10 5 IF(B)15,15,20 15 . . . .

produces a code involving T 00 twice. The second T 00 is eliminated.

(2) It eliminates the code RECALL N, ST 00 when not needed, such as occurs in the translation

D05J=1,N

- (3) It replaces the code RE 00, . . ., T 00, by T 00, . . . provided that the user has not specified a COMPILE R for 00 or that the meaning has not changed. This item is for situations in which 00 is used as a temporary variable prior to a zero being needed in 00 for any reason.
- (h) A final check is made to eliminate all recalls of nondimensional variables that are in the window. If an RE R is directly preceded by PRINT, format, or INDIR, ST R, or any jump, SEARCH, Ll, or combinations of these, and the next previous operation is a  $\rho R$ , then the RE R is eliminated. Similarly, if the code is STORE, X or RECALL, Y followed by the above nondestructive steps, followed by RECALL, X, the last two steps are eliminated. After any jump, a STOP or RETURN is allowed in place of SEARCH.

## 6. EXAMPLE

This lengthy example illustrates the use and effectiveness of this compiler. The example consists of a main program that calls the subroutine POLRT, taken without change from the 360 Scientific Subroutine Package. 5 The POLRT is used to find the complex roots of

<sup>&</sup>lt;sup>5</sup>System/360 Scientific Subroutine Package, Version III, Programmer's Manual, Program Number 360A-CM-03X, ed. GH20-0205-4, IBM Corporation, White Plains, NY (August 1970), 181-183.

real polynomials; the calling program must supply the degree M and the M+1 coefficients XCOF of the polynomial. A simple main program was therefore written to read these inputs, call POLRT, and write the output. Figure 5 shows the program, with comments to give motivations for each of the COMPILE statements that affect the translation. The MAIN was made a subroutine to allow the arguments to be aligned identically to POLRT. This is a common maneuver to save steps when subroutines are called, but it does require the COMPILE S statement. The dimensions of arrays were guessed at to try to allow large-degree polynomials in a fully expanded Wang 520.

In the \*FTC card for MAIN, no options were specified; hence, both the source code and the translation output (name table, translation, and reference table) are printed. Figures 6 to 8 show the translation output. These have been previously described, except that the reference table (fig. 8) now contains three additional lines: The top work register is 09 (registers 10 to 15 contain the arguments); only single DO loops are involved (because of reading and writing arrays); and one is the maximum number of work registers needed to translate any statement.

In the \*FTC card preceding POLRT, the AUTO option was specified. Since POLRT calls no other routines and all of its nonargument variables are initialized at every entry, many steps are saved by allowing the compiler to assign basic registers for the most frequently used variable. If some of its nonargument variables were not initialized every time, or if POLRT had been called in a loop, then some register conflicts could arise. Then AUTO (nn-mm) would have to be used by the programmer to specify a more limited range. It is always up to the programmer to insure that no conflicts arise when many routines are compiled together.

Figure 9 contains the source listing of POLRT (the original comments have been deleted); figure 10, the name table; figure 11, the translation; and figure 12, the reference table. In figure 12, the work registers are 09, 08, and 07, and register 01 is used for the single DO loops. If AUTO were not specified, registers 02-06 would never be used. The name table in figure 10 assigns N, X, Y, SUMSQ, and IFIT to these registers. In addition, the compiler senses that TEMP can be assigned to one of the work registers (07), since it is a temporary variable in a single domain that does not use register 07. Similarly, the variable DX was assigned register 01, because it is defined and used in a single domain outside a DO loop, and the variables XTZ, L, and ITEMP were all assigned register 00. (When L is used as an index, it is assigned

<sup>&</sup>lt;sup>1</sup>H. Bloom and A. Hausner, FORTRAN IV Compiler for the Wang 520 Calculator, Harry Diamond Laboratories TM-73-15 (July 1973).

```
SUBROUTINE MAIN(XCOF, COF, M, ROOTR, ROOTI, IER)
      DIMENSION XCOF(25), COF(25), ROOTR(24), ROOTI(24)
      READ() M
COMPILE W M(1500-1-3)
C THIS W CARD PRINTS M ON THE READ IN FORMAT 1500 WITH SPACES BEFORE
 AND AFTER THE PRINT.
     N = M + 1
      READ() (XCOF(I), I=1,N)
COMPILE W XCOF(0510,1,0)
C. THIS W CARD PRINTS THE COEFFICIENTS XCOF ON THE READ IN FORMAT 0510.
    CALL POLRT(XCOF, COF, M, ROOTR, ROOTI, IER)
COMPILE S POLRT(E, E, E, E, E, E), 0-0
C THIS S CARD INDICATES THAT NO ARGUMENTS OF MAIN NEED BE SAVED OR
C RESTORED, AND THAT NO ARGUMENTS OF POLRT NEED BE INPUTTED OR OUT-
C PUTTED. THE REASON MAIN WAS MADE A SUBROUTINE WAS TO MATCH ARGUMENTS
C FOR THIS PURPOSE.
      IF(IER.NE.O) GOTO 5
      WRITE() (ROOTR(I), ROOTI(I), I=1,M)
COMPILE W ROOTR(0010,0,1)
COMPILE W ROOTI (1110.0.0)
C. THESE W CARDS PRINT THE REAL AND IMAGINARY ROOTS IN FORMATS 0010 AND
 1110, SEPARATING EACH PAIR OF ROOTS BY SPACES.
      STOP
    5 WRITE() IER
COMPILE W IER (0700.0.1) .
C THIS W CARD PRINTS THE VALUE OF IER IN THE FORMAT 0700 AFTER A SPACE.
      STOP
      END.
```

Figure 5. Program illustrating use of COMPILE statements.

# NAME TABLE

INDEX	NAME	DIM	TYPE	ARG
			4	
1	MAIN	0	1	50
2	XCDF	25	0	15
3	COF	25	0	14
4	M	0	- 1	13
5	ROOTR	24	0	12
6	ROOTI	24	0	11
7	IER	0	1	10
8	N	0	1	20
9	I	. 0	1	20
10	POLRT	C	, Э	50

Figure 6. Name table for program MAIN.

register Ol.) These temporary variables would have been assigned auxiliary registers even if AUTO had not been specified. The net effect of AUTO was just the assignment of variables in registers 02-06.

When the reader examines the translation of POLRT in figure 11, he must keep in mind some of the new optimization features automatically used. For example, N+1 is a common subexpression in the line below statement number 35, NXX=N+1 and two lines following KJl=N+1. Register 00 is used for this subexpression, computed in steps 56 to 58, and later recalled in step 64. Also, KJl is stored in register 07 (step 65) for use in the DO loop following (steps 72 and 88). Register 07 was available for that local domain and one step was saved by this action. (If KJl had been a basic register, one step would have been wasted. This feature is implemented before register assignments and does sometimes lead to wasted steps.)

Following the POLRT source deck was an \*LDR card and an \*END card with no options listed. The storage maps (fig. 13) and program listing (fig. 14) are printed regardless of any options that affect only entry point codes, mark numbers, and registers assigned. In figure 13, the arrays were assigned registers 16 to 113 as listed under LIMITS in the storage map for MAIN. (These are not repeated in the storage map for POLRT, because the dimensions there were given as one. All variable dimensioned arrays should have dimension one.) The other columns are almost self explanatory. The INDEX, NAME, DIM, TYPE, and ARG columns are repeated from the name table (fig. 6) only for variables that must be assigned space. Columns for COM and EQU are for variables in common or equivalence. The LOC column contains the location of the variable or, in the case of an array, the location of the pointer. locations can be used to recall these variables with the indirect code. To recall variables with the direct code, the REG column lists the direct register name. The LIMITS column is applicable only to arrays.

```
0 0900 SHIFT MARK 45 0804 J IF 0
1 MAIN 46 0800 SEARCH
2 0802 PRINT 47 5
3 0015 CLRDSP 48 0001 E 1
4 0004 E 4 49 0601 ST 01
5 0609 ST 09 50 0900 SHIFT MARK
2 0802 PRINT 47 5
3 0015 CLRDSP 48 0001 E1
4 0004 E 4 49 0601 ST 01
5 0609 ST 09 50 0900 SHIFT MARK
6 0903 SHIFT STOP 51 LABD1
7 0613 ST 13 52 0802 PRINT
8 0802 PRINT 53 0015 CLRDSP
9 1500 SP-RE 00 54 0701 RE 01
10 0802 PRINT 55 0609 ST 09
11 0015 CLRDSP 56 0712 RE 12
12 0609 ST 09 57 0209 + 09
13 0001 E 1 58 1511 INDIR
14 0209 + 09 59 0709 RE 09
15 0901 SHIFT STORE 60 0802 PRINT
16 N 61 0010 E .
17 0001 E 1 62 0701 RE 01
18 0601 ST 01 63 0609 ST 09
19 0900 SHIFT MARK 64 0711 RE 01
18 0601 ST 01 63 0609 ST 09
19 0900 SHIFT MARK 64 0711 RE 01
20 14800 65 0209 + 09
21 0509 ST 09 66 1511 INDIR
22 0902 SHIFT MARK 64 0711 RE 11
20 14800 65 0209 PRINT
24 0600 ST 00 69 1110 SHFT F 10
25 0715 RE 15 70 0713 RE 13
26 0209 + 09 71 0603 ST 00
27 0302 E 2 72 0001 E 1
28 0200 + 09 71 0600 ST 00
29 0903 SHIFT STOP 74 0902 SHIFT ALPHA
30 1511 INDIR
31 0609 ST 09 76 0805 SIN
31 0609 ST 09 76 0805 SIN
31 0609 ST 09 76 0800 SHIFT STOP
32 0802 PRINT
33 0510 / 10 78 0903 SHIFT STOP
34 0801 RECALL 79 0903 SHIFT STOP
35 0806 SIN
36 0600 ST 00 81 0802 PRINT
37 0001 E 1 82 0015 CLRDSP
38 0201 + 01 83 0710 RE 10
39 0902 SHIFT ALPHA 84 0802 PRINT
40 0806 SIN 85 0700 RE 00
41 0806 SIN 85 0700 RE 00
41 0800 SIN 85 0700 RE 00
41 0801 RECALL 86 0903 SHIFT STOP
44 0806 SIN 85 0700 RE 00
45 PRINT STOP 87 0914 SHIFT END
           44 0710 RE 10
```

Figure 7. Wang relative code for program MAIN.

# REFERENCE TABLE

```
MAIN
N
16 35
POLRT
47 80
LABOO 20 42
LABOI 51 77

THE FIRST AVAILABLE WORK REGISTER IS 9
MAX DO LOOP NEST IS 1
MAX REG RANGE IS 1
```

Figure 8. Reference table for program MAIN.

FORTRAN COMPILER POLRT AUTO

```
SUBROUTINE POLRTIXCOF.COF, M, ROOTR, ROOTI, (ER)
DIMENSION XCDF(1), COF(1), ROOTR(1), ROOTI(1)
  DOUBLE PRECISION XO, YO, X, Y, XPR, YPR, UX, UY, V, YT, XT, U, XT2, YT2, SUMSQ,
 1 DX. DY. TEMP. ALPHA
   IFIT=0
N=M
IER=0
   IF (XCOF(N+1))10, 25, 10
10 IFIN 15, 15, 32
15 IER=1
20 RETURN
25 IER=4
GQ TO 20
$0 1ER=2
GO TO 20
32 IF(N-36) 35,35,30
 NX=N+1
NX=N+1
N2=1
35 NX=N
 % KJ1 * N+1
00 40 L=1,KJ1
MT=KJ1-L+1
40 COFINTI=XCOFILI
45 X0=.00500101
Y0=0.01000101
Y0=0.01000101
IN=0
50 X= X0
X0=-10.0*Y0
  AD=-10.0+X
```

Figure 9. Program listing for POLRT.

```
X= XO
    Y=Y0 ...
    IN = IN + 1
    GO TO 59
 55 IFIT=1
    XPR=X
    YPR=Y
 59 · ICT=0
 60 UX=0.0
    UY=0.0
    V =0.0
    YT=0.0
    XT=1.0
    U=COF(N+1)
   IF(U) 65,130,65
 65 DO 70 I=1, N
    L = N - I + 1
    TEMP=COF(L)
    XT2=X*XT-Y*YT
    YT2= X*YT+Y*XT
    U=U+TEMP*XT2
    V=V+TEMP*YT2
    FI=I
    UX=UX+FI*XT*TEMP
    UY=UY-FI*YT*TEMP
   XT=XT2
 70 YT=YT2
    SUMS Q=UX+UX+UY+UY
    IF (SUMSQ) 75,110,75
 75 DX=(V*UY-U*UX)/SUMSQ
    X = X + DX
    DY =- (U*UY+V*UX )/SUMSQ
    Y = Y + DY
 78 IF(DABS(DY)+DABS(DX)-1.00-05) 100,80.80
 80 ICT=ICT+1
    IFIICT-5001 60,85,85
 85 IF(IFIT)100,90,100
 90 IFLIN-51 50,95,95
95 IER=3
    GO TO 20
100 DO 105 L=1,NXX
    MT=KJ1-L+1
    TEMP=XCOF(MT)
    XCOF(MT)=COF(L)
105 COF(L)=TEMP
    ITEMP=N
   N=NX.
    NX= ITEMP
F(IFIT) 120,55,120
110 IF(IFIT) 115,50,115
115 X=XPR ...
    Y=YPR
120 IFIT=0
122 IF(DABS(Y)-1.00-4*DABS(X)) 135,125,125
```

Figure 9. Program listing for POLRT (cont'd).

```
125 AL PHA= X+ X
    SUMS Q= X * X + Y * Y
    N=N-2
    GO TO 140
130 X=0.0
    NX = NX - 1
    NXX = NXX - 1
135 Y=0.0
    SUMS Q=Q.Q
    ALPHA=X
    N = N - 1
140 COF(2)=COF(2)+ALPHA*CUF(1)
145 DO 150 L=2.N
150 COF(L+1)=COF(L+1)+ALPHA*COF(L)-SUMSO*COF(L-1)
155 ROOT [[N2]=Y
    ROOTR(N2) = X
    N2 = N2 + 1
    IF (SUMSQ) 160, 165, 160
160 Y=-Y
    SUMSQ=0.0
    GO TO 155
165 IF(N) 20,20,45
    END
```

Figure 9. Program listing for POLRT (cont'd).

Listings under LABEL and MARK indicate the mark code given for each label. In figure 13 MAIN used marks 0000 to 0002, and POLRT used marks 0003 to 0107.

No conflicting marks or registers exist, despite that register 114 was assigned to the variable N in MAIN and the variable XO in POLRT. This is part of an optimization scheme that equivalences variables in several programs. In effect, the variables are placed in common. To be a candidate for such treatment, the domain of a variable must not contain any function or subroutine calls, and the variable must not be in an EQUIVALENCE or COMMON statement. Such variables are considered local to the problem and can be destroyed after leaving the program.

The final listing is shown in figure 14. Array pointers are first set up (steps 02 to 13). Next, the routines MAIN and POLRT are listed with the assigned registers, labels, and entry points replacing the symbolic ones. All statement references in figures 8 and 12 are so replaced. Finally, the large numbers .500101D-02 and .1000101D-01 needed in steps 99 and 103 of POLRT (fig. 11, 12) are generated as data in the program code in steps 720 to 735. Because any steps converted to

NAME TABL	E
-----------	---

INDEX	NAME	MID	TYPE	ARG
		1.		
	POLRT	0	o	50
2	XCOF	1	0	15
2	COF	ī	0	14
4	М		1	13
5	ROOTR	1	9	12
6	ROOTI	1	0	1.1
7	TER : XO YO	0 1 0 0 0 0	1	10
8	XO	0	0	20
9	YO	0	U	20
10	X	0	0	3
12	XPR	0	0	20
12	YPR	0	0	20
14	UX	0	0	20
15	UY	0	Ö	20
16	V	0	0	20
1.7	YT	0	0	20
18	XT	0	0	20
19	U	0	0	20
20	XT2	0	0	0
21	YT2	0	0	20
22	SUMSQ	0	0	
24	DX	0	0	20
25	TEMP	0	0	20
26	ALPHA	O	.0	20
27	LEIT	0	i	6
28	N	0	1	2
11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34	NX	0	1	20
30	NXX.	0	1	20
31	N2	0	1	20
32	KJ1	0	1	20
33	L	0	1.	0
3.4	MT	0.	ી. જાજ	20
35 36 37 38	KJ1 L MT IN		1	50 15 14 13 12 11 10 20 20 20 20 20 20 20 20 20 20 20 20 20
37	ICT	0	1	20
38	FI	0	0	20
39	ITEMP	0	1	- 0
39			-	0,

Figure 10. Name table for POLRT.

Figure 11. Wang code for POLRT.

```
164 0901 SHIFT STORE 219 0608 ST 08
```

Figure 11. Wang code for POLRT (cont'd).

STEP	CODE BUTTON	KEY	STEP	CODE BUTTO	N KEY	
220	0801	RECALL	275	ΥT		
	and the same of th				DECAI	ī
2.2.2	2620 V	0.0	277	3831	RECAL	L
222	3438 X	08 09 STORE	2//	Y 1 Z		
223	3239 +	39	2/8	3901 SHIF	T STORE	
224	3901 SHIFT	STORE	279	ΥT		
225	Y T 2	O7 O9 O0 O9 RECALL O9 STORE	280	0702 R	E 02	
226	0707 RE	07	281	0600 S	T OO	
227	3639 ST	29	282	0001	E 1	
228	O7OO RE	<b>၁</b> ၁	283	0201	+ 01	
229	3439 X	39	284	0902 SHIE	T ALPHA	Δ
230	0801	RECALL	285	0806	SIN	
231	11	NCC PCC	296	2822	CEADO	L
232	1210	20	200	3033	SEARC	,П
222	0001 CHIEF	67605	201	001	0.5044	
233	0401 2HILL	O7 O9 RECALL O9 RECALL	288	J801	RECAL	L
234	U		289	UX		
235	0707 RE	07	290	0812	X = *2	
236	3639 ST	09	291	0605 S	T 05	
237	0801	RECALL	292	0801	RECAL	L
238	YT2		293	UY		
239	0409 X	39	294	0812	<pre><pre></pre></pre>	
241	2821	RECALL	295	1205	± 05	
261	V	REGALE	304	0004 5415	TINE	0
242	0300	22	290	0904 3016	CCADO	
242	0209 +	RECALL  OP STORE  OI STORE  OP RECALL  OP O7 OP RECALL	291	0800	SEARC	H
243	0401 2H1F1	STURE	298	110		
244	V		299	C801	RECAL	.L
245	2721 RE	01	300	V		
246	1931 SHIFT	STORE	301	3631 S	T Ol	
247	FI		302	0801	RECAL	L
248	3609 ST	39	303	UY		
249	1080	RECALL	304	0401	X 01	
250	ΧT		305	0801	RECAL	1
251	0409 X	39	306	11	0	
252	1717 DE	27	307	2608 6	T 00	
252	3/3/ KL	20	300	0000	DECAL	
253	2001	25	300	6901	KELAL	
254	3801	RECALL	309	UX		
255	UX		310	0408	X 08	
256	0209 +	09	311	0301	- 31	
257	3931 SHIFT	09 STORE	312	0705 · R	E 05	
258	UX		313	0501	/ 01	
259	0801	RECALL	314	0203	+ 03	
260	UY			0801	RECAL	L
	0509 ST	09		U		
	0801	RECALL			T 09	
	FI		318		RECAL	ı.
		08		UY		
	3831	RECALL		0409	X 09	
	YT	RECALL				
		0.0	321		RECAL	L
		08		V	T 00	
	0707 RE				T 08	
	0408 X			0801	RECAL	L.
270	0309 -		325			
271	0901 SHIFT	STORE			X 08	
	UY		327	0209	+ 39	
273	0700 RE	00		2012	CHS	
	0901 SHIFT				T 09	

Figure 11. Wang code for POLRT (cont'd).

Figure 11. Wang code for POLRT (cont'd).

439 0800 SEARCH

440	55		495	0901 S	HIFT	STORE
441	0800 120 0900 SHIF	SEARCH	496	NX		
442	120		49.7	3831		RECALL
443	0900 SHIF	L WALK	498	NXX		
444	110		499	0509	ST	09
445	0705 R 0904 SHJF	- 06	500	3731		E l
445	0904 SHIF	TJNFO	501	3439	-	319
447	0900	SEAFCH				
448	50		503	NKK		
449	0801 XPR	RECALL	504	0900 S	HIFT	MARK
45)	XPR		505	135		
451	0603 5	T 03	506	0104	Ŧ	34
452	108C	RECALL	507	01 05	T	35
453	Abb		508	0703	RE	
454	0604 S 0900 SHIF	T 04	509	0901 S	HIFT	STORE
455	0900 SHIF	L MARK	510	, ALP	HA	
450	120		511	0001		E 1
457	0106	7 06	512	3332	-	25
458	0106 0704 RI	04	513	0900 S	HIFT	MARK
459	0913 SHIF	r ABS		140		
460	76 79 S	7 )9	515	0002		E 2
401	31133	1111	516	0609 0714	ST	09
462	0703 RI	0.3	517	3714	RE	14
463	2913 SHIF	ABS	518	0209	+	09
464	0902 SHIFT 1104 SHFT 1 0309 0012	ALPHA	519	2001		E 1
465	L104 SHET	- 24	520	0608	ST	08
456	3309	- 09	521	7714	RE	14
467	2012	CHS	522	0208	+	08
<b>400</b>	0702 31111	wit Lind	7 2 3	とフレレ		TANTE
469	0805 0800	J IF +	524	0708	RE	08
477	0800	SEAFCH	525	3638	ST	3 8 C
471	135 0703 R		525	1080		RECALL
472	0703 R	03	527	0408 1511	НΑ	
	0609 S		529	0408	`	08
474	0209	19				
475	3931 SHIF	STORE	530	0209	+	09
			531	3732		E 2 "
477	0703 RI 0812 0605 ST 0704 RF	03	532	0601	ST	21
478	0812	X**2.	533	0900 S	HIFT	MARK
479	0605 5	0.5	534	D 03		
483	3734 RI	. 04	535	3731	RE	01
481	0812	x * * 2	536	3639	ST	29
		05	537	0001		F 1
	0002	3 2	533		+	<b>3</b> 9
434	0302 -	. 02	539	0714	RE	14
485	0800	SEARCH	540	72 79	+	39
486	140		541	1070	BF	01
	0900 SHIFT	MARK	542	0608	ST	08
	130		543	2714	RE	14
		03	544	0208	+	38
	0801	RECALL	545	1511		INDIR
491	NX		546	0708	RE	0.8
	0609 \$1	09	547	2638	ST	
	0001	F 1	548	3831		RECALL
494	0309 -	. 09	549	ALPI	HΑ	

Figure 11. Wang code for POLRT (cont'd).

```
605 0704 RE 04
606 0012 CHS
607 0604 ST 04
608 0105 T 05
609 0800 SEARCH
610 155
611 0900 SHIFT MARK
612 165
613 0100 T 00
614 0702 RE 02
615 0902 SHIFT ALPHA
516 0806 SIN
617 0800 SEARCH
618 20
619 0800 SEARCH
620 45
621 0914 SHIFT END
569 0600 | ST 00
 570 0001 E 1
571 0201 + 01
 572 0902 SHIFT ALPHA
 573 0305 SIN
574 0800 SEAR
                         SEARCH
 575 D03
 576 0900 SHIFT MARK
 577 155
 578 3831
                         RECALL
 579 N2
 580 0609 ST 09
581 0711 RE 11
582 0209 + 09
+ 09

583 0704 RE 04

584 1511 INDIR

585 0609 ST 09

586 0801
                     RECALL
 587 N2
 587
588
0609
581
09
589
0712
RE
12
590
0209
09
591
0703
RE
03
592
1511
INDIR
593
C609
ST
09
594
0801
RECALI
                     RECALL
 595 N2
 596 0609 ST 09
 597 0001 E 1
598 0209 + 09
 599 0901 SHIFT STORE
 600 N2
 601 0705 RE 05
 602 0904 SHIFT J NE 0
 603 0800 SEARCH
 604 165
```

Figure 11. Wang code for POLRT (cont'd).

# REFERENCE TABLE

POLRT					1					
_XO			16.6		-101	1112	124			
Y0 -					105	124	130		V 11 " "	
XPR'					150	450				
YPR					153	453			•	
UX .				1 2 X	163	255	258	289	309	325
ÜY .					165	260	272	293	303	319
V					167	241	244		322	
-Y.T				3-1	169		216.	266	279	
XT				7. 4	172	205	221	250	275	-=-
U=				6-1	182	231	234	306	316	
YT2					225	238	277	300	250	
DY			11.11.		333	336				
ALPHA				5.7	476	510	527	549		
NX				• 2	55	432	436	491	496	
NXX					60	421	498	503	12.70	
N2					63	579	587	595	600	
KJI				• .• .	67	393	201	272	000	
IN	• •	• • •		• •	1.08	1.36	141	373	0.00	
ICT					158	353	358	3.13		
FI			<b>4</b> 5., •	. H .	247		370			
25				4	17	30	1 1 5			
32					23	42				
20	• • •			•	27	: 34	40	385	618	
30			1		36		- 40	303	0.10	
D00 -					71	95		-8 -		
45			• • •	• •	97	620				
50					110	381	448	*		
59				• •	143	155	440			
55					145	443				
60				2.5	160	367				
130					185	488				
D01	• • •	• • • •		• .•	189	287				
110					298	444	. 2.	"		
100		• • • •		• •	35l		. 387		-	
D02	• •	• • •		• •	391		301			
120					442	428		4-		
135				38.18.						
140				1 1	471	505				
D03		7			486	514				
155				•	534 577	575				
165					604	.610				
103	5001010	200000	-02:			612				
		000000		• •		03				
-4,	190019	. 000000	, - 01			03				

THE FIRST AVAILABLE WORK REGISTER IS 9 MAX DO LOOP NEST IS 1 MAX REG RANGE IS 3

Figure 12. Reference table for POLRT.

LOADER MAIN					
	STORAGE M	IAP .			
INDEX NAME DIM	TYPE	ARG COM	EQU LOC	REG	LIMITS
2 XCOF 25	0	15 0	0 15	UP15	16 - 40
3 COF 25		14 0	0 14	UP14	41 - 65
5 ROOTR 24	0	13 0	0 12	UP13 UP12	66 - 89
6 ROOTI 24	0	11 0	0 11	UP11 UP10	90 - 113
8 N 0		20 . 0	0 114	RE02	0 - 0
			*		

0 - 0

LABEL MARK 5 LABOO LABO1

GADER POLRT

ting.		STORAGE	MAP	ر بره می و نهرستند و با		lament and a supplement		72
INDEX	NAME	DIM TYPE	ARG	CON	EQU	LOC	REG LIMITS	
2	XCOF	1 0	15	ō	0	15	UP15 0 -	<b>翻</b> (
3	COF	1 0	14	0	0	14	UP14 0 -	1
4	M	0 1	13	0.	0	13	UP13 0 -	(
5	ROOTR	1 0	12	0	0	12	UP12 0 -	<b>AND C</b>
6	ROOTI	0 1	11	0	. 0	11	UP11 .3 -	
7	LER	0 1	1.0	0	0	10	UP10 0 -	300
. 8	XO	0 0	20	0	0	114	RE02 0 -	380
9	YO	0 0	20	0	0	115	RE03 0 -	强(
10	X	0 . 0	3	0	0	3	UP03 0 -	型(
11	Y	0 0	4	0	0	4	UP04 0 -	(2000年)
12	XPR	0 0	20	- 0	0	116	RE04 0 -	<b>1</b>
13	YPR	0 0	20	0	0	117	RE05 0 -	遷(
14	UX	0 0	20	0	0	118	RE06 0 -	· · · · (
15	UY	0 0	20	0	0	119	RE07 0 -	3 (
16	٧	0.0	20	0	0	120	RE08 0 -	. P
17	YT	Ó O	20	0	0	121	RE09 0 -	0
18	XT	00	20	. 0	0	122	RE10 0 -	0
19	U	0 0	. 20	0	0	123	RE11 0 -	. 0
20	XT2	0 0	0	0	0	0	UP00 0 -	0
21	YT2	0 0	20	0	0	124	RE12 0 -	<b>1</b> (
22	SUMSQ	0 0	. 5	0	0	- 5	UP05 0 -	, a c
23		0 0	1	0	0	1	UP01 0 -	0
24	DY	0 0	20	0	0	125	RE13 0 -	0
25	TEMP	0 0.	7	0	0	7	UP07 0~-6	0
26	ALPHA	0 0	20	0	0	126	RE14 0 -	- TO
27	IFIT	0 1	6	0	0	6	UP06 0 -	3 C
28	N	0 1	2	0	0	2	UP02 0 -	0
29	NX	0 1	20	0	0	127	RE15 0 -	0
	NXX	0 1	20	0	. 0	128	SP-UPOQ 0 -1	3 0

Figure 13. Storage map for MAIN and POLRT.

1 N2 0 1 20 0 129 SP-UPDT	0	400 -	0
32 KJ1 0 1 20 0 0 130 SP-UPO2	0	-	0
33 L 0 0 0 0 UP00	0		. 0
35 IN 3.0 1 20 0 0 131 SP-UP03	. 0	-	0
36 1CT 0 1 20 0 0 132 SP-UPO4	0	-	0
38 FIG. 37 SP-UP05	0	-	0
39 ITEMP 0 0 0 UP00	0	<del>-</del>	0
- [1] 가는 사용하는 경기 등을 통해 보는 이 기계를 보고 하는 것이 되는 것이 되는 것이 되었다. 그는 것이 되는 것이 되는 것이 되었다. 			1,574.14

			9.1	
LAB	EL	1	MAR	K
7	V 107	- 1	14	41.41
25				2
	The State of		10. 87 300	
32	Q1 14		22	4
20-	100		P. 10. 1	5
30	17	3, 5		6
DOO		le le	. B.L.	7
45		13.61		8
50		1		0
59		1	40.7	0
11 mm / 7g		v ing .	3.7	0
55	13			1
60		300	1	2
130	4.		1	3
DOL			1	4
110			1	
100	Trick.	1. 1	10	0
D02	ments and	1 14	2-7	
			10	
120		10.0	10	-
135	12		10	3
140			10	4
003	1		10	5
155		1	10	
165			10	

Figure 13. Storage map for MAIN and POLRT (cont'd).

0	3900	SHIFT	MAPK	55	LOCI	F(X)	01
i	1000	F(X)	00	56	2713	RE	10
2	0001		F 1	57	2804		J IF D
3	0005		E 5	58	0900		SEARCH
4	0615	ST	15	59	0001	ALL UP	01
- 5	JJJ4		F 4	50	0001		E 1
6	ეეეე		E 0	61	J501	ST	01
7	7614	ST	14	62	0900	SHIFT	MARK
Ω	2006		F 6.	63	2000	ALI UP	02
9	1005		E 5	64	0802		PRINT
10	1512	ST	12	55	2015		CLRDSP
11	0003		F B	66	0701	RE	01
12	2003		E S	67	0509	ST	0.9
1.3	7611	ST	11	63	7712	S E	12
14	7872		PRINT	69	7233	+	9
15	0015		CERUSE	70	1511		INDIP
16	0004		F 4	71-	0709	RE	Có
1.7	0609	ST	39	7.2	0802		PRINT
1.8	0903	SHIEL	STOP	73	0010		Ε.
[0	0513	ST	13	74	2721		71
50	0832	60.66	PRINT	75	3639	\$1	39
21	1500	SP-RF	00	76	7711	Q <u>C</u>	11
22	0802		bt Iv.L	7.7	0200	+	09
23	0015	СТ	CLPDSP	78	1511	0.0	INDIR
24	0609	21	00	79	2779	RE	09
25 26	2220		F 1	9 D 9 I	0802	SHET F	TRINT
27	0209	SHIFT	STURE	95 21	0713	SELLE	13
23	3732	- SHIFT	02	33	2673	ST	00
29	0001	IV C	E 1	94	2221	3 1	E 1
30	3501	S 1	01	35	32.01	+	01
31	3331	SHIFT	MARK	36	390?	SHIFT	
32	0000	ALL UP	20	87	0205	511611	SIN
33	2629	ST	)9	89	0900		SEARCH
34	1)91)2	SHIFT	ALPHA	99	0003	ALL UP	02
35	11-03	SHET F	13	90	0903	SHIFT	
36	1610	ST	77	91	0900	SHIFT	MARK
37	3715	RE	15	35	1000	ALL UP	01
3.8	0209	+	0.9	93	J302		PRINT
39	0002		E 2	94	1015		CLRDSF
47	2533	+	00	15	3713	₹ [	1 3
41	0903	SHIFT	ctub	36	0002		BSINI
42	1511		INCIR	9.7	0700	RE	70
4.3	0609	ST	()9	38	22.13	SHIFT	STOP
44	16.15		PRINT	99	ככיכ	SHIFT	MARK
4.5	1513	/	10		1001	, F ( ×)	
46	1831		ACCALL	101	3106	T	06
47	7702		0.2	102	0713	RE	1.3
4.7	0500	ST	C.O.	103	0608	SI	
49	))))]		FI	134	0602	ST	35
51		+ 	)]	105	0110	T	10
51	0902	SHIFT	VERHA	105	0708	P.F.	0.8
52	0806		SIN	107	2609	ST	79
= 53	2802		SFARCH	109	10001		F 1
54	JUJU	ALL UP	UC.	109	1509	+	0.7

Figure 14. Final Wang code listing.

Figure 14. Final Wang code listing (cont'd).

220	0703	ΝF	03	275	0503	+	09
221	1407	X	0.7	276	0714	RE	1 4
	0901	SHIFT	STORE	277	0209	+	29
223	0702	RF	02	278	1511		INDIR
			_			25	
224	0700	RE	33	279	0709	RE	
225	0601	ST	O L	280	3931	SHIFT	STORE
226	0708	RE	08	281	2711	RE	1.1
227	2421	Х	01	282	0904	SHIFT	J NE O
228	2921	SHIFT	STOPE	283	0800	3	SEARCH
						411 110	
229	3733	PF	33	284		ALL UP	13
230	0727	ЬE	07	235	3331		E 1
231	7633	ST	73	286	0601	ST	01
232	2721	RF	21	287	3300	SHIFT	MARK
233	0604	ST	24	288	2014	ALL UP	1.4
234	0801		RECALL	289	0702	RE	C2
235		ŞP					
	0803		03	290	C60C		0.0
236	3679	ST	09	291	0701	RE	_
237	3001		F 1	292	0300	_	
239	0203	141	39	293	0001		E 1
239	1000	SHIFT	STORE		0200	+	00
240	0803	SP	03		0609		09
			SEARCH				
241	0800				0714	RF	14
242		ALL UP	10	277		+	9
243	0000	SHIFT	MARK	298	1511		INDIR
244	1100	ALL LIP	1.1	299	0709	RE	09
245	0001		E 1	300	0607	ST	07
246	1616	S.T.	26	301	0703	RE	03
247	2723	RF	03	30?	3633		00
	7901						
248			STURE	303	2801		RECALL
249	0704	RE	04	374	2713	RE	10
250	2704	PF	04	3 3 5	0400	X	30
251	0901	SHIFT	STOPE	306	7734	RE	24
252	0705	RF	0.5	337	3638	ST	38
253	0900	SHIFT	MARK	308			RECALL
254	0010		10	309		RE	09
		466 0-					
255	0000		E 0	310		X	
255	0901		STOPE	311	0300	~	
257	0904	SP	04	312	3733	R E	33
258	0900	SHIFT	MAPK	313	2609	ST	09
259	J012	ALL UP	12	314	1080		RECALL
260	2000		FO	315	2739	RE	
261	2001			316		X	
262							
	7736	PE	06	317		RE	
263	7971	SHIFT	STORE	313	0608	ST	38
254	0707	9.4	07	319	1080		RECALL
265	0901	SHIFT	STURE	320	0710	RE	10
266	3738	B.E.	28	321	0408	X	08
267	0901	SHIFT	STORE	322	0209	141	39
268	7779	F.E.	39	323	3901	SHIFT	STORE
		7.5	E 1				
269	)))]			324	0712	RE	12
270	3931	SHIFT	STUDE	325	0707	RF	0.7
271	0710	BE	10	326	26.39	ST	29
272	0702	SE	32	327	7777	9.E	70
273	26 19	ST	79	328	0409	X	09
274	3001	_	E 1	329	1680	,	RECALL
	1001			2 1. 7	3001		HECHLE

Figure 14. Final Wang code listing (cont'd).

STEP	CODE	BUTTON	KEY	STEP	CODE	BUTI	ron	KEY
330	0711	RE	1.1	206	0200			CEADCU
		+	00	386	0014	ALI	110	SEARCH 14
	0901	CHIET	STORE	207	3001	ALL	UP	RECALL
	3711	31111	11	300	2706		0.5	
	0707		07	300	0013		KE	06 X**2
			22	307	0605		СТ	A T T C
	0801	31	DECALL	390	2021			RECALL
	0712		RECALL	202	0707		0.5	RECALL
	3439	X	20	392	0101		KE	7442
	3831	^	12 09 RECALL	304	0205			05
347	0778		18	205	0207	CHI	ET	J NE O
341	0730		OB O9 STORE	306	1911	3111	1 1	SEARCH
342	0901	SHIET	STORE	207	2015	A I. I	HD	15
343	0708	PE	08	308	0801	MLL	Ur	RECALL
344	0701	P.C.	0.1	300	0708		DE	VECHEE
345	0901	SHIFT	STORE	400	0601		CT	01
346	0805	SP	08 01 STORE 05 09	421	0801		31	RECALL
347	2629	ST	29	402	0707		DE	0.7
348	2821	3 ,	RECALL	403	0401		Y	01
	0710		10	404	0801		^	RECALL
	0409		22	4.05	0711		DE	1.1
			27	406	0608		ST	08
352	0409	X	09	407	0801		J 1	RECALL
353	3831	X	RECALL	408	0706		RE	06
354	2726	RF	26	409	0408		X	08
355	0209	+	RECALL 36 09 STORE 06 RECALL	410	0301		_	01
355	0901	SHIFT	STORE	411	2725		RE	35
357	0706	RE	06	412	3531		1	31
358	0801		RECALL	413	0203		+	03
359	0707	RE	07 09	414	0801			RECALL
360	0609	ST	29	415	0711		RE	11
361	0801		RECALL	416	3639		ST	39
362	0805	SP	05	417	0801			RECALL
363	0608	ST	08	418	0707		RE	07
364	0801		RECALL	419	0409		X	09
365	0709	RE	09	420	0801			RECALL
366	0408	X	08	421	0708		RE	08
367	0707	RE	07	422	0608		ST	08
368	0408	X	07 09 RECALL 05 08 RECALL 09 08 07 08 07 08	423	1080			RECALL
369	0309	_	09	424	0706		RE	06
370	0901	SHIFT	STORE	425	0408		X	08
	0707				2239			39
	0700				2100			CHS
		SHIFT			3639		ST	
	0710 0801				0705		RE	
			RECALL	430	0509	CUI	-	
	0712 0901		12 STORE	431	0901 0713	211	RE	STORE
	0709				3234		KE +	04
	0702	RE			0801		*	RECALL
	0600	ST			0713		RE	
	0001	31	E 1		0913			ABS
	0201	+			0609		ST	
		SHIFT			0701		RE	
	0806	0	SIN		3913			ABS
					3 . 2 3	5111	. *	

Figure 14. Final Wang code listing (cont'd).

STEP	CODE	BUTTON	KEY	STEP	CUDE	BUTTON	KEY
447	1209	+	9	495	0300	-	00
441	0100	Ţ	0.0	496	0001		E 1
442	0001		5 1	497	3200	+	0.0
443	0902	SHIFT	ALPHA	498	06 0 9	ST	39
444	1105	SHET F	05	499	7715	RE	15
445	0309	-			05 0 3	+	
446	0012		CHS	501			INDIR
447		SHIFT	ALPHA	502		RF	
448	0905		J IF +				
	0800		SEAFCH		2700	RF	
450		Ţ	-13	505	0609		
451	0801		PFCALL	506	0715	RE	
452		SP	J 9	101	0207	+	
	3639	ST	59	500	0/01	R E S T	31 38
	0001		F 1	510	0716	R E	
455 456	0209 0901	+ T = 1 t = 2	STORE	511	0608 0714 0208	+	
	0804	31111	0.6	2 6 6	1511		INDIR
	0005	. SP	E 5		3738		
	0902		ALPHA			,,,,	INDIR
	1002	F(X)	0.2	515	3639		09
	0309		09	516	3737		27
462				517	1511		INDIR
463	2090	SHIFT	CHS ALPHA	518		ST	08
	0805		J IF +	519	1080		RECALL
465	0300		SEARCH	520	0800	SP	20
466	0012	ALL UP		521	0600	ST	00
467	0706	RE		255	0001		F 1
	0804		J IF O	523	9201	+	01
	0800			524	0902	SHIFT	ALPHA
473	7100	Ţ	DD RECALL	525	CR 06		NIS
471	2821			りとい	0800	Ψ.	SEARCH
472	0803	SP		527	)702	T RE	
473			) 9 E 5	620	0600	C T	00
474	0005 0309	_	36	530	03,01	31	RFCALL
475			CHS	531	0715		15
	0902	SHIFT	ALPHA				25
478			J IF +				00
	0800		SEARCH	534	3901	SHIFT	STORE
		ALL UP	SEARCH 09	535	0715	RE	15
481	)003		F 3	535	7705	RE	05
	0610	ST	10		0904		J NE O
	0800		SEARCH		0800		SEARCH
494		ALL UD	15	539	J011	ALL UP	11
485	0900	SHIFT	MUDK	540	0300		SFARCH
486	0110	T	0.0	541		1	32
437	0001		E L			SHIFT	
498	7671	ST	21		0015		15
489	7917	SHIFT	MAPK	544		RE	06
493	2121	Ţ	01	545		SHIFT	J NE O
491	0801	e n	SECULE	540 547		A 1 1 1 1 1 1 A	SEARCH 09
432	2622	SP ST	72	543		ALL JP	RECALL
493	0701	S I RF	)) )]		0704	RE	
774	OIOT	P. F	OF	277	0104	1× L	J 1

Figure 14. Final Wang code listing (cont'd).

STEP	CODE	BUTION	KEY	STEP	CODE	BUTTON	KEY
550	0603	ST	O3 RECALL	635	3134	Т	04
551	0801		RECALL	636	0105	T	05
552	0705	RE	35	607	0703	RE	
553	0604	ST	04	608	0901	SHIFT	
554	0933	SHIFT	MARK	609	0714	RE	14
555	0102	Ţ	02	610	3331		E 1
556	0106	1	06	611	3332	-	02
557	37 14	KE	J4	612	0900	SHIFT	
550	0413	20111	0.0	616	0104	Т	E 2
560	0100	31	00	615	0602	5.7	39
561	0703	PE	0.3	616	0714	0.5	14
562	2913	SHIFT	ARS	617	02.09	+	09
563	0902	SHIFT	AL PHA	518	0001	•	E 1
564	1104	SHET E	04	519	0608	ST	08
565	0309	_	09	620	0714	RE	14
566	0012		CHS	621	0208	+	08
567	3972	SHIFT	RECALL 35 04 MARK 02 06 34 ABS 39 00 03 ABS 4LPH4 09 CHS ALPH4	622	1511		INDIR
568	0805		J IF +	623	0708	RE	08
	0800		SEARCH	624	0608	ST	08
	0103	T	03	625	0801		RECALL
571	0703	RE	03	626	2714	RE	
	0609	ST	39	627	0408	Х	08
573	0209	+	39	628	1511		INDIR
	0901	SHIFT	STURE	529	0209	+	09
575	0714	K E	14	630	0002		E 2
575 577	0703	FE	U 3	631	7071	SUIET	31
	0605	с т	SEARCH 03 03 09 09 STORE 14 03 X**2 05 04 X**2 05 E 2 02 SEARCH	632	0105	SHIFT	
579		RF	04	634	0701	R.F.	01
580	3812		X**2	635	2629	ST	09
581	0205	+	25	636	0001		E 1
582	2003		E 2	637	0209	+	09
583	0302	-	02	638	2714	RE	
584	0800		SEARCH	639	2239	+	29
535	01.04	T	04	640	3731	RE	01
596	0900	SHIFT	MARK	541	0608	ST	08
	0013	ALL UP	13	542	0714	RE	14
588	0103	T		643	3238	+	
	0801	٥.۶		644	151-1	0.6	INDIR
591	0609	KE	15 09	645	0708	RE	08
-	0009	31	E 1			ST	
593	0301	_	09		0801	DE	RECALL 14
594	0901	SHIFT	STORE	649		X	38
595	0715	PE	15	550	1511	^	INDIR
596	1080		RECALL	651	0709	RE	09
597	0800	SP	00	652		+	08
598	0609	ST	09	653	0701		01
599	1000		E 1	654	2627	ST	37
	0309	_	09	655	0001		E 1
571	0911	SHIFT	STORE	656	0307	-	07
572		SP	))	657		RE	14
	3933	SHIFT	MARK	658	0207	+	07
504	0103	T	03	659	1511		INDIR

Figure 14. Final Wang code listing (cont'd).

```
660 0707
                             RF 27
                                                                            710 0900 SHIFT MARK
                                                                 710 0900 SHIFT MAPK
711 0107 T 07
712 0100 T 00
713 0702 PE 02
714 0902 SHIFT ALPHA
715 0806 SIN
716 0800 SEARCH
717 0005 ALL UP 05
718 0800 SEARCH
719 0006 ALL UP 08
720 0300 T 00
721 0100 T 00
722 0000 E 0
724 0001 E 1
725 0001 E 1
725 0000 E 0
727 0500 / 00
728 0200 + 00
729 0100 T 00
730 0000 E 0
731 0000 E 0
732 0100 T 00
733 0100 T 00
734 0000 E 0
735 0100 T 00
736 0914 SHIFT END
660 0707 RE 07

661 0607 ST 07

662 0705 RE 05

563 0407 X 07

664 0308 - 08

665 1511 INDIR

666 0609 ST 09

667 0702 RE 02

668 0600 ST 00
                                                                            711 2127 T 27
                               - 08
INDIR
669 0001 E 1
670 0201 + 01
 671 0902 SHIFT ALPHA
672 0806 SIN
673 0800 SEAPCH
674 0105 T 05
 675 0900 SHIFT MARK
 575 0106 T C5
677 0801 PECALL
678 0801 SP 01
679 0609 ST 09
680 J711 PF 11
681 0209 + 09
682 0704 RE 04
683 1511 INDIR
684 0609 ST 09
685 0301 SP 01
686 0801 SP 01
687 0609 ST 09
688 0712 FF 12
639 0209 + 09
690 0703 RE 03
                                 PECALL
 677 0801
                                                                            736 3914 SHIFT END
690 0703 RE 03
691 1511 INDIR
692 0609 ST 09
693 0801 RECAL
                                RECALL
694 0801 SP 01
695 0609 ST 09
696 1001 E 1
697 0209 + 09
 698 3901 SHIFT STORE
699 0801 SP 01
700 0705 RF 05
 701 0904 SHIFT J NF 0
                       SEARCH
T 37
 702 0800
 703 0107
704 0704 RE 04
705 0012 CHS
706 0604 ST 04
                             T 05
SEAFCH
 707 0105
  708 0800
                                T 06
  709 0106
```

Figure 14. Final Wang code listing (cont'd).

register use must start with a step number divisible by 8, unused steps preceding the data are replaced with the GO code (0830). Thus, there could be as many as seven GO statements preceding the large number data registers (the example in fig. 14 has none). These could be replaced by a small user-added subroutine or simply ignored.

The final printed output (fig. 15) contains miscellaneous information that may be useful to a user of the program. The verify program (VP) number is given for the user to check when the program is laoded in the machine. Marks and registers used are specified so the user can determine what is left for other purposes. The total number of steps (program and registers) is given for the same purpose. If this number is greater than 1848, the program is too large for the machine. Finally, the entry points are given to facilitate any editing the user may wish to perform.

ENTRY POINTS NAME MARK

MAIN 1000 POLRT 1001

THE VERIFY PROGRAM NUMBER IS 7832

REGISTERS USED

16 - 133

MARKS USED

0 - 107 1000 - 1001

118 REGISTERS WERE USED

1681 STEPS OF TOTAL STORAGE WERE USED

Figure 15. Entry point and register information.

The punched cards provided as output by the compiler are put into a local program to punch the actual Wang cards read by the card reader (sect. 2). The first card shows the hexadecimal number of cards following, and each succeeding card contains 40 steps at two columns per step in hexadecimal representation. (The code 0512 is written as 5C; i.e., high-order and low-order code each takes a column.)

To use the program once the code is in the machine, it is necessary to know how to control the I/O flow. Hence, it is wise to write the instructions on the back of the final Wang input cards. For this program, such instructions would probably look like the following:

Purpose: To find roots of real polynomial  $\sum_{i=0}^{M} C_{i+1} x^{i}$ , 0 < M  $\leq$  25

(Translation of POLRT)

(Explanation)

VP = 7719--load only at step 0000

Put Printer ON.

1. Key f0.

Entry point for MAIN

- 4 appears in display. Enter degree
   M of polynomial and key GO (printed 4 is index of M (fig. 6)
   M)
- 3. 2.i appears in display. Enter C 2 is index of XCOF and key GO (printed C). (fig. 6)

After last coefficient is entered, roots are computed and printed (X real, I imaginary), unless they are not found. Then, error code is printed (labelled E)

E = 1 means M is less than 1.

E=2 should never arise since we have made

E = 2 means M is greater than 36.

provision for M < 26

E = 3 means unable to determine root with 500 iterations on 5 starting values.

 $E = 4 \text{ means } C_{M+1} = 0.$ 

Other information could be indicated if it might be useful (we use the symbol " $\rightarrow$ " to mean "is stored in register"):

 $M \rightarrow 13$ 

Coefficients → 16 to 40

Real part of roots → 66 to 89

Corresponding imaginary part of roots → 90 to 113

Marks used: f00-1, 0000-T07

Total steps used: 1680 (736 for code at low end and 944 at high end)

With these instructions, only the source and object listings need be saved if trouble can be anticipated.

It is hoped that the detail of this example helps clarify how to use the compiler. Experiments will probably clarify the use further and suggest other techniques for improvements. The authors would appreciate comments or suggestions for this purpose.

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### APPENDIX A. -- ADAPTABILITY OF SYSTEM TO OTHER COMPUTERS

The compiler system has been exclusively written in FORTRAN IV and is independent of the word length (32 bits) or character code (EBCDIC), except for the following cases:

- (a) FUNCTION DIG--This function takes a digit character code (left justified) and right justifies as numerical digit.
  - (b) FUNCTION NUMBER--This function is the opposite of DIG.

The program takes approximately 280k bytes of core on the IBM 370/195. In addition, two routines reference a direct access file (logical unit 1) that contains records 466 words in length. Approximately 130 records can be stored on the disk. The two routines are listed below:

- (a) GET (A, I, O) -- Gets array A from the Ith record on the file.
- (b) PUT (A, I, O) -- Puts array A into the Ith record on the file.

The routine BIN takes the loader codes (n Wang steps) and punches out an object deck containing 40 steps to a card, each step represented as a two-digit hexadecimal code. A leader card gives the count of the number of cards (also in hexadecimal) in the object deck. This object deck is then converted into the Wang card form through a conversion routine written for the IBM 1130. This routine could be altered to produce the Wang object deck directly.

A listing of the compiler is available upon request. The system contains approximately 11,000 cards, including 4000 comment cards. The source code may be obtained by sending a tape to the authors. The code will be placed on a nine-track tape at 800 BPI on one file. Also available upon request is a listing of the IBM 1130 program for converting the object decks.

In converting the system to a smaller computer, it will be necessary to overlay the system. Fortunately, the system is divided into six autonomous subsystems (sect. 3 of the main body of the report) with the important tables being passed through labelled commons or disk files.

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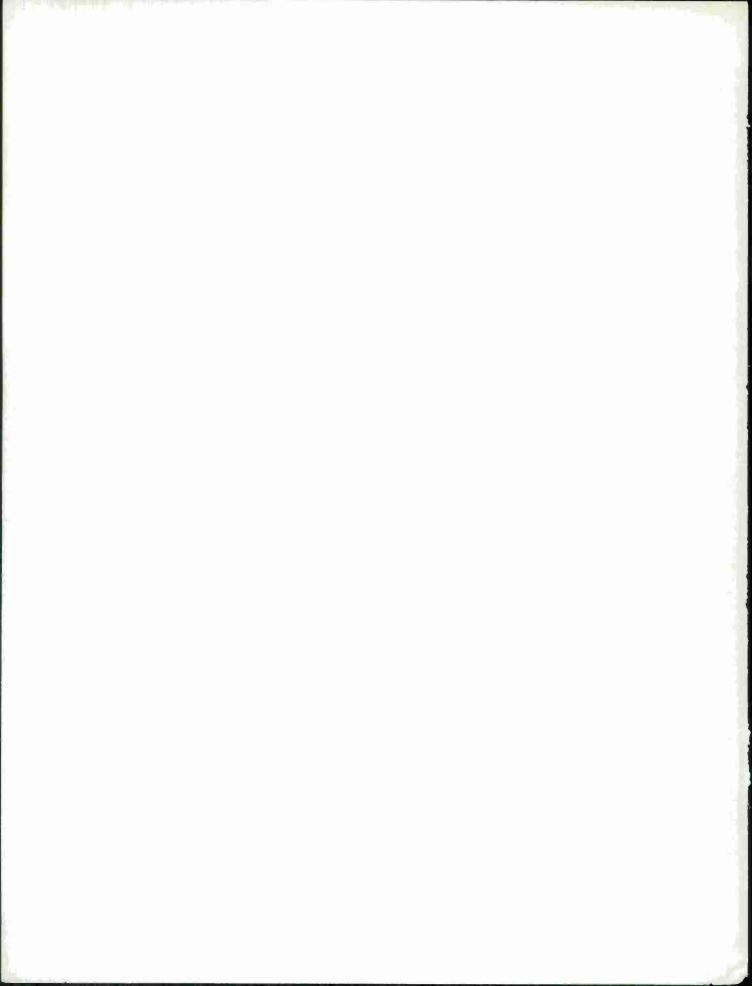
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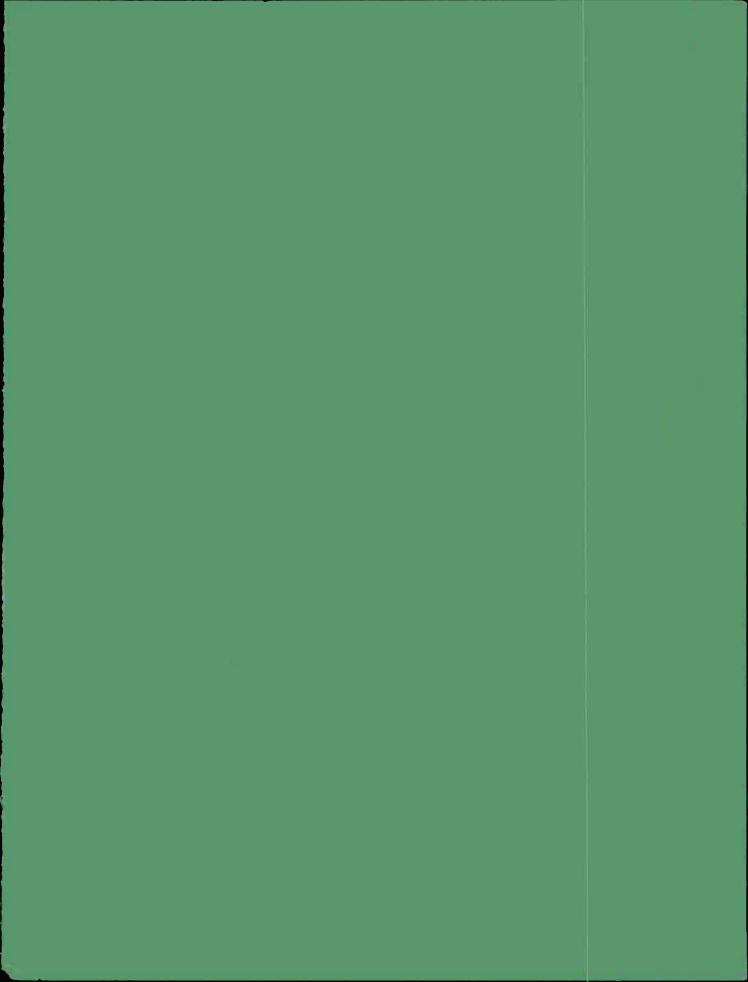
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